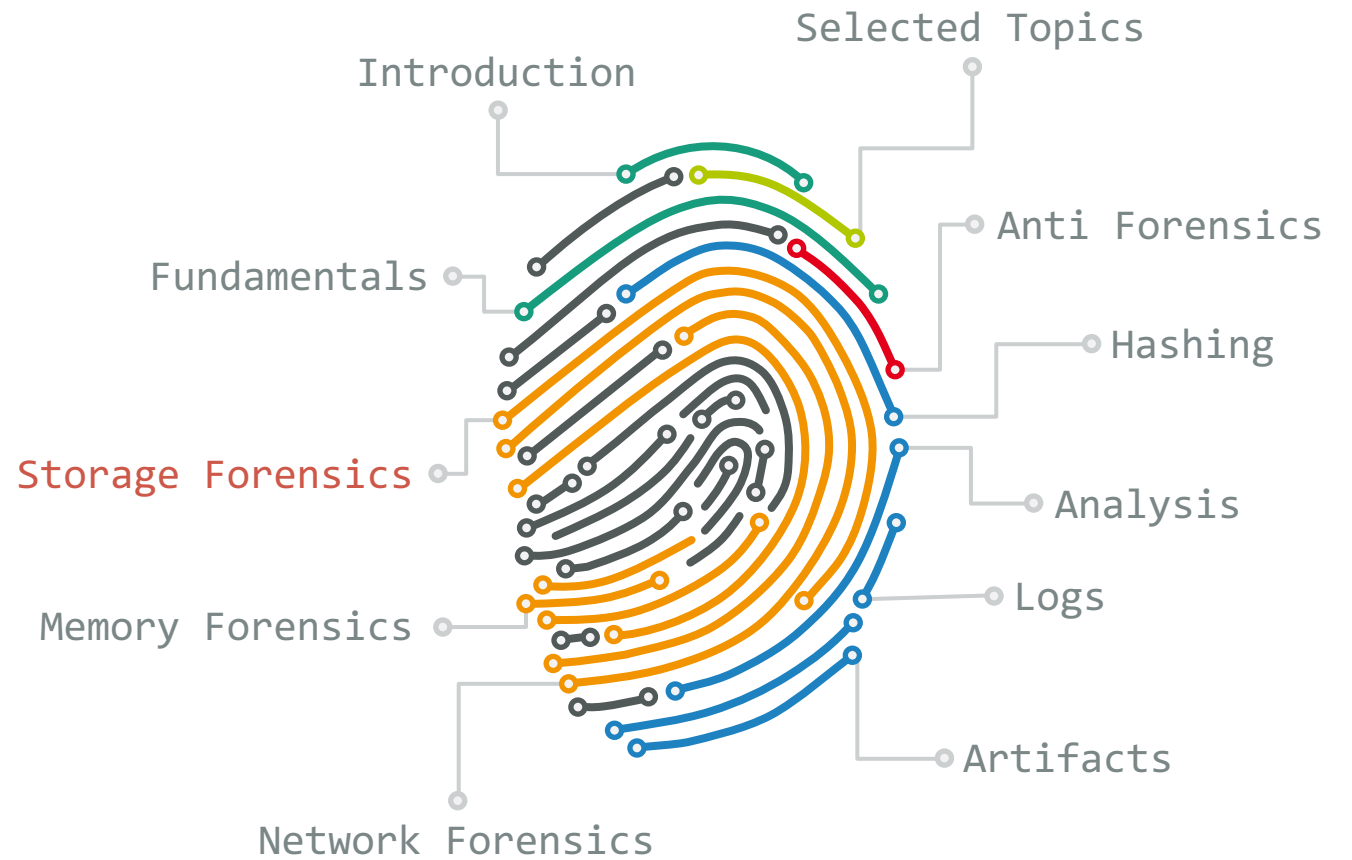


Prof. Dr. Elmar Padilla et al.

Digitale Forensik

02 - Storage Forensics

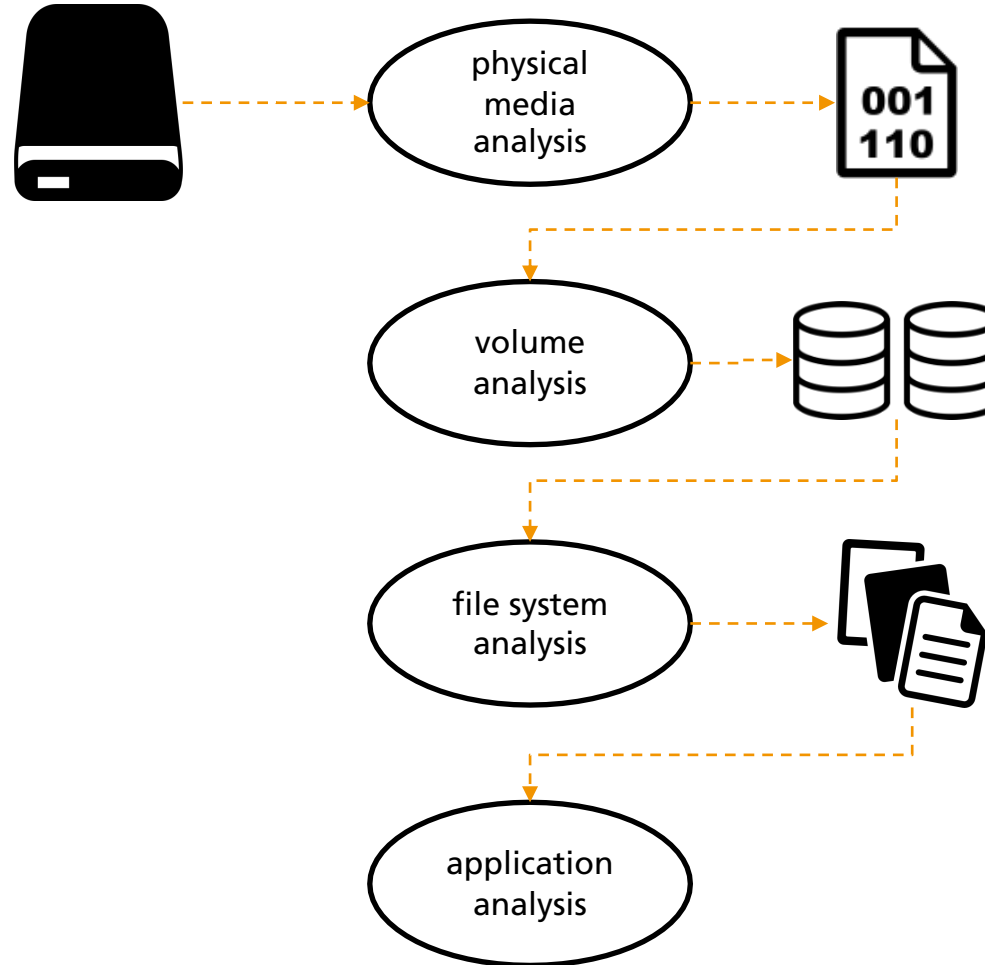




File System Forensic Analysis



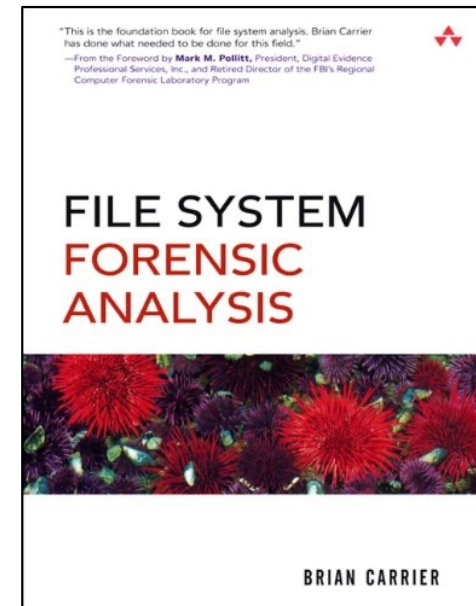
Dr. Brian Carrier
Basis Technology



Sleuth Kit



Autopsy

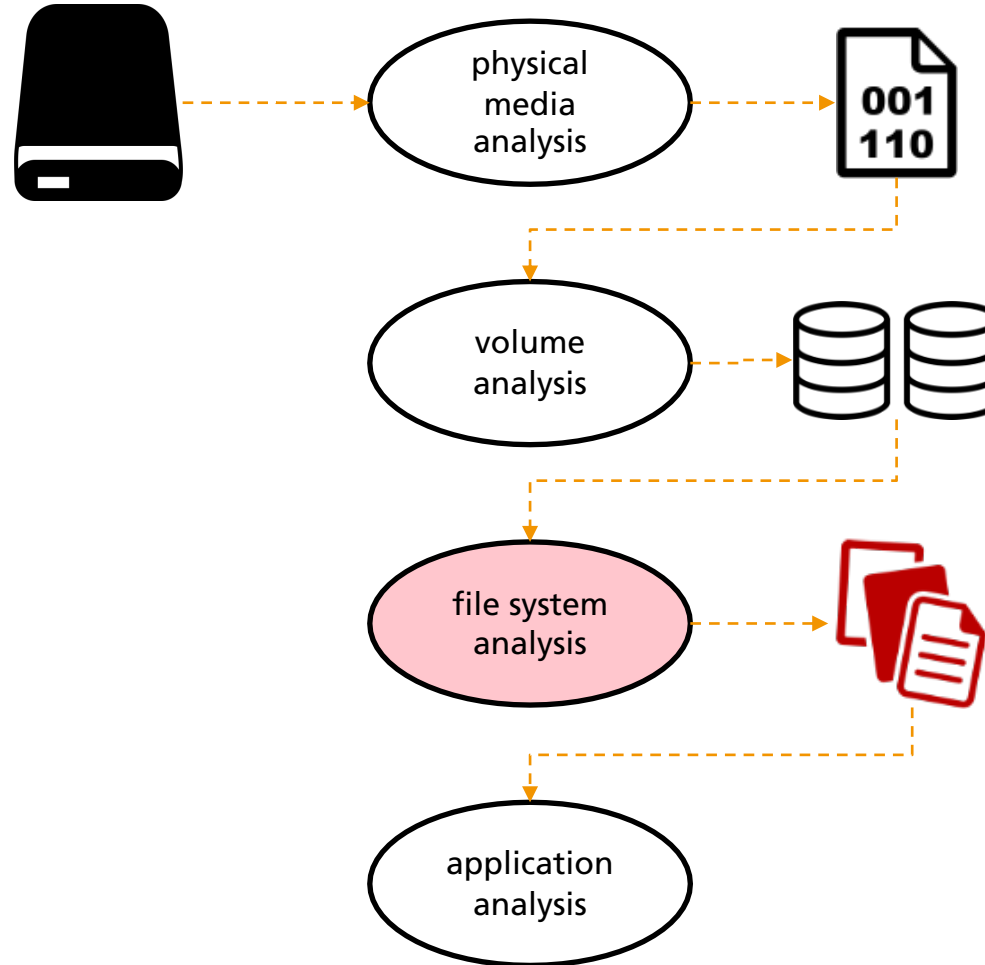




File System Forensic Analysis



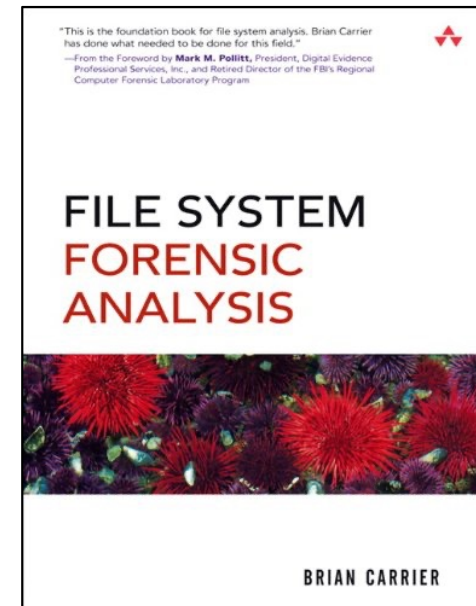
Dr. Brian Carrier
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Sleuth Kit

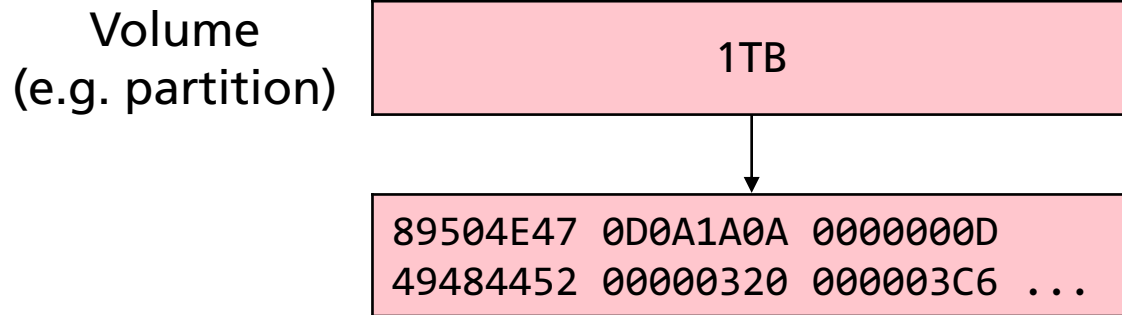


Autopsy



File System Analysis

Why File Systems?



Storing data directly on a volume works, but is infeasible for multiple files

1. How to **reference** a file?
2. How to quickly find the **start** of a file and collect **metadata** (such as timestamps)?
3. How to organize **available space**?
4. How to implement features such as **logging**?
5. How to **organize** the aforementioned concepts?



image.dd

```
89504E47 0D0A1A0A 0000000D 49484452 00000320
000003C6 08060000 001B6DFE F7000000 01735247
4200AECE 1CE90000 00786558 49664D4D 002A0000
00080004 011A0005 00000001 0000003E 011B0005
00000001 00000046 01280003 00000001 00020000
87690004 00000001 0000004E 00000000 00000090
00000001 00000090 00000001 0003A001 00030000
00010001 0000A002 00040000 00010000 0320A003
00040000 00010000 03C60000 0000B58D 47E10000
00097048 59730000 16250000 16250149 5224F000
00400049 696d6167 6530312e 6a7067D4 2455792E
AB5A88C6 2427829A 93758C03 338D3F47 D7D1ACA5
E255FEA7 D144136F 048C0283 0B9DC623 B2B29418
40D1DC83 D224F1DC 08842162 7E04951E 24FCCD44
63927382 F43080C6 242B8A37 1A4D0CF7 7DBABFB7
A7BEEABD DFBDAB6A 5777757F CFCBDADF AE7AFF77
```

File Systems

„The motivation behind a file system is fairly simple: computers need a method for the long-term storage and retrieval of data. File systems provide a mechanism for users to store data in a hierarchy of files and directories. A file system consists of structural and user data that are organized such that the computer knows where to find them. In most cases, the file system is independent from any specific computer.“

Dr. Brian Carrier

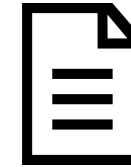


image.dd

```

89504E47 0D0A1A0A 0000000D 49484452 00000320
000003C6 08060000 001B6DFE F7000000 01735247
4200AECE 1CE90000 00786558 49664D4D 002A0000
00080004 011A0005 00000001 0000003E 011B0005
00000001 00000046 01280003 00000001 00020000
87690004 00000001 0000004E 00000000 00000090
00000001 00000090 00000001 0003A001 00030000
00010001 0000A002 00040000 00010000 0320A003
00040000 00010000 03C60000 0000B58D 47E10000
00097048 59730000 16250000 16250149 5224F000
00400049 696d6167 6530312e 6a7067D4 2455792E
AB5A88C6 2427829A 93758C03 338D3F47 D7D1ACA5
E255FEA7 D144136F 048C0283 0B9DC623 B2B29418
40D1DC83 D224F1DC 08842162 7E04951E 24FCCD44
63927382 F43080C6 242B8A37 1A4D0CF7 7DBABFB7
A7BEEABD DFBDAB6A 5777757F CFCBDADF AE7AFFF7
  
```

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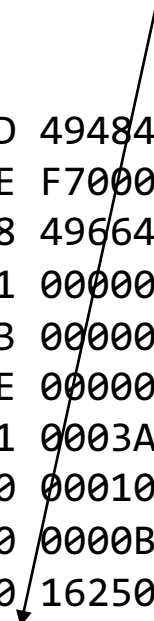
Dr. Brian Carrier



image.dd

File name

89504E47	0D0A1A0A	0000000D	49484452	00000320
000003C6	08060000	001B6DFE	F7000000	01735247
4200AECE	1CE90000	00786558	49664D4D	002A0000
00080004	011A0005	00000001	0000003E	011B0005
00000001	00000046	01280003	00000001	00020000
87690004	00000001	0000004E	00000000	00000090
00000001	00000090	00000001	0003A001	00030000
00010001	0000A002	00040000	00010000	0320A003
00040000	00010000	03C60000	0000B58D	47E10000
00097048	59730000	16250000	16250149	5224F000
00400049	696d6167	6530312e	6a7067D4	2455792E
AB5A88C6	2427829A	93758C03	338D3F47	D7D1ACA5
E255FEA7	D144136F	048C0283	0B9DC623	B2B29418
40D1DC83	D224F1DC	08842162	7E04951E	24FCCD44
63927382	F43080C6	242B8A37	1A4D0CF7	7DBABFB7
A7BEEABD	DFBDAB6A	5777757F	CFCBDADF	AE7AFF7



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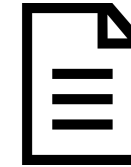


image.dd

File content

89504E47	0D0A1A0A	0000000D	49484452	00000320
000003C6	08060000	001B6DFE	F7000000	01735247
4200AECE	1CE90000	00786558	49664D4D	002A0000
00080004	011A0005	00000001	0000003E	011B0005
00000001	00000046	01280003	00000001	00020000
87690004	00000001	0000004E	00000000	00000090
00000001	00000090	00000001	0003A001	00030000
00010001	0000A002	00040000	00010000	0320A003
00040000	00010000	03C60000	0000B58D	47E10000
00097048	59730000	16250000	16250149	5224F000
00400049	696d6167	6530312e	6a7067D4	2455792E
AB5A88C6	2427829A	93758C03	338D3F47	D7D1ACA5
E255FEA7	D144136F	048C0283	0B9DC623	B2B29418
40D1DC83	D224F1DC	08842162	7E04951E	24FCCD44
63927382	F43080C6	242B8A37	1A4D0CF7	7DBABFB7
A7BEEABD	DFBDAB6A	5777757F	CFCBDADF	AE7AFFF7

File Systems

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Operating systems need to be able to parse and understand all of this file system information!

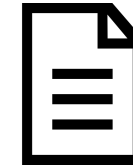
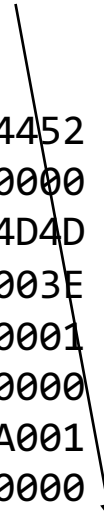


image.dd

Timestamp

89504E47	0D0A1A0A	0000000D	49484452	00000320
000003C6	08060000	001B6DFE	F7000000	01735247
4200AECE	1CE90000	00786558	49664D4D	002A0000
00080004	011A0005	00000001	0000003E	011B0005
00000001	00000046	01280003	00000001	00020000
87690004	00000001	0000004E	00000000	00000090
00000001	00000090	00000001	0003A001	00030000
00010001	0000A002	00040000	00010000	0320A003
00040000	00010000	03C60000	0000B58D	47E10000
00097048	59730000	16250000	16250149	5224F000
00400049	696d6167	6530312e	6a7067D4	2455792E
AB5A88C6	2427829A	93758C03	338D3F47	D7D1ACA5
E255FEA7	D144136F	048C0283	0B9DC623	B2B29418
40D1DC83	D224F1DC	08842162	7E04951E	24FCCD44
63927382	F43080C6	242B8A37	1A4D0CF7	7DBABFB7
A7BEEABD	DFBDAB6A	5777757F	CFCBDADF	AE7AFFF7



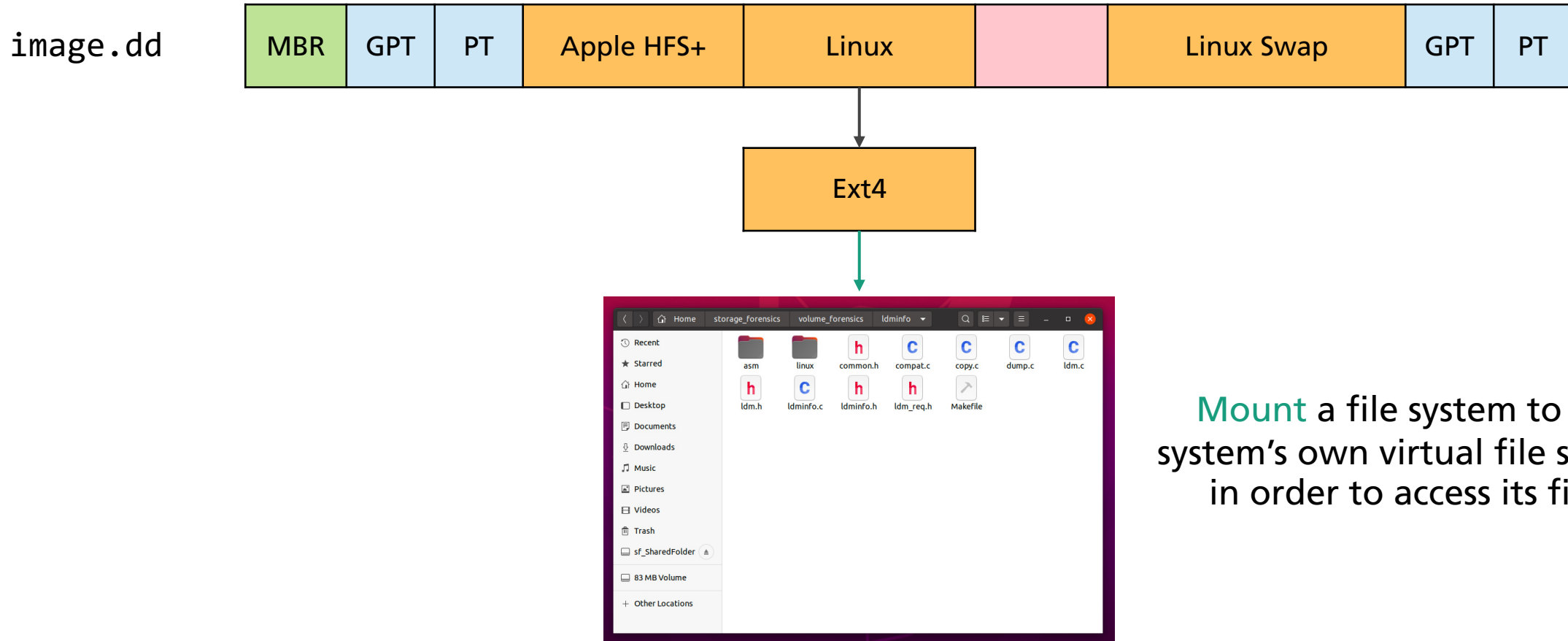
File Systems

The screenshot shows a GitHub repository view for the 'linux/fs/fat' directory. At the top, there is a breadcrumb navigation 'linux / fs / fat /' and a 'Go to file' button. Below this, a commit by 'torvalds' is shown, with a merge tag 'linux-kselftest-kunit-5.15-rc1'. The commit message is 'Merge tag 'linux-kselftest-kunit-5.15-rc1' of git://git.kernel.org/pu...'. The commit hash is 'c815f04' and it was made 'on 2 Sep'. A 'History' link is also present. Below the commit information, a list of files is displayed, each with a file icon, the filename, a brief commit message, and the time since the last commit.

File	Commit Message	Time
..		
.kunitconfig	fat: Add KUnit tests for checksums and timestamps	2 months ago
Kconfig	fat: Add KUnit tests for checksums and timestamps	2 months ago
Makefile	fat: Add KUnit tests for checksums and timestamps	2 months ago
cache.c	fat: new inline functions to determine the FAT variant (32, 16 or 12)	3 years ago
dir.c	treewide: Remove uninitialized_var() usage	15 months ago
fat.h	fs: make helpers idmap mount aware	9 months ago
fat_test.c	fat: Add KUnit tests for checksums and timestamps	2 months ago
fatent.c	block: move the bdi from the request_queue to the gendisk	2 months ago
file.c	Merge tag 'idmapped-mounts-v5.12' of git://git.kernel.org/pub/scm/lin...	8 months ago
inode.c	mm: require ->set_page_dirty to be explicitly wired up	3 months ago
misc.c	fat: Add KUnit tests for checksums and timestamps	2 months ago
namei_msdos.c	fat: handle idmapped mounts	9 months ago
namei_vfat.c	fat: handle idmapped mounts	9 months ago
nfs.c	treewide: Replace GPLv2 boilerplate/reference with SPDX - rule 282	2 years ago

<https://github.com/torvalds/linux/tree/master/fs>

File System Analysis



Mount a file system to the system's own virtual file system in order to access its files.

File System Analysis

mount

mount(8) - Linux man page

Name

mount - mount a filesystem

Synopsis

```
mount [-lhV]
```

```
mount -a [-fFnrsvw] [-t vfstype] [-O optlist]
```

```
mount [-fnrsvw] [-o option[,option]...] device|dir
```

```
mount [-fnrsvw] [-t vfstype] [-o options] device dir
```

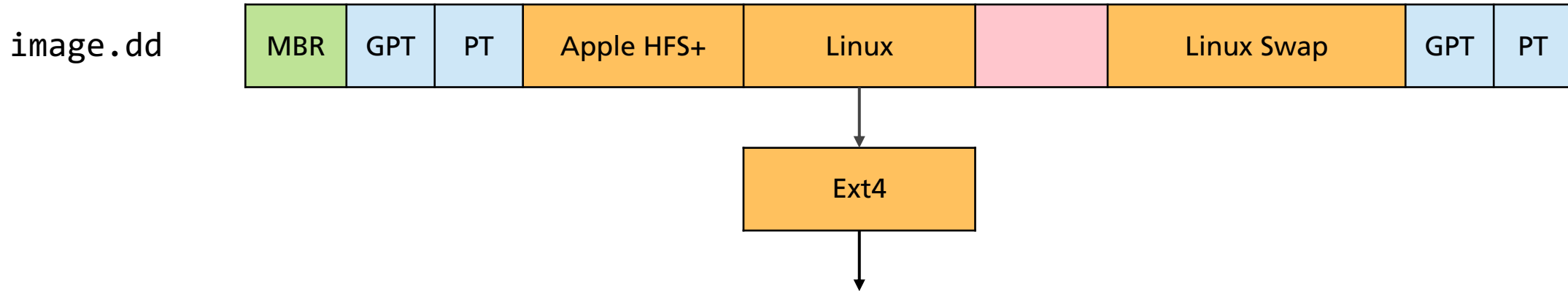
Description

All files accessible in a Unix system are arranged in one big tree, the file hierarchy, rooted at /. These files spread out over several devices. The **mount** command serves to attach the system found on some device to the big file tree. Conversely, the **umount(8)** will detach it again.



Mount has a **-o ro** (for **read-only**) option!

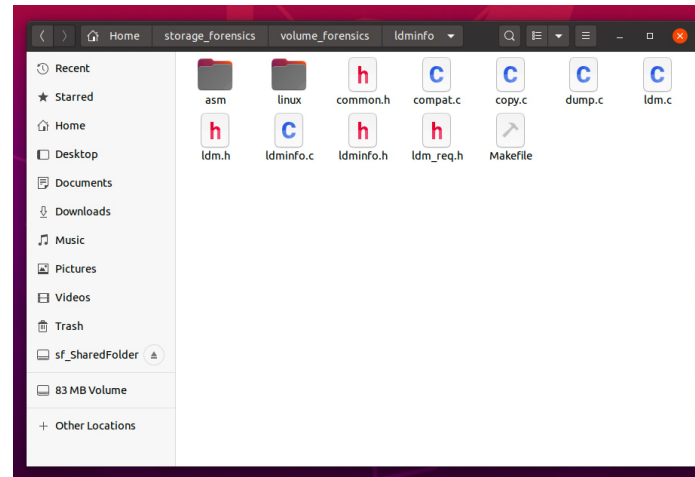
File System Analysis



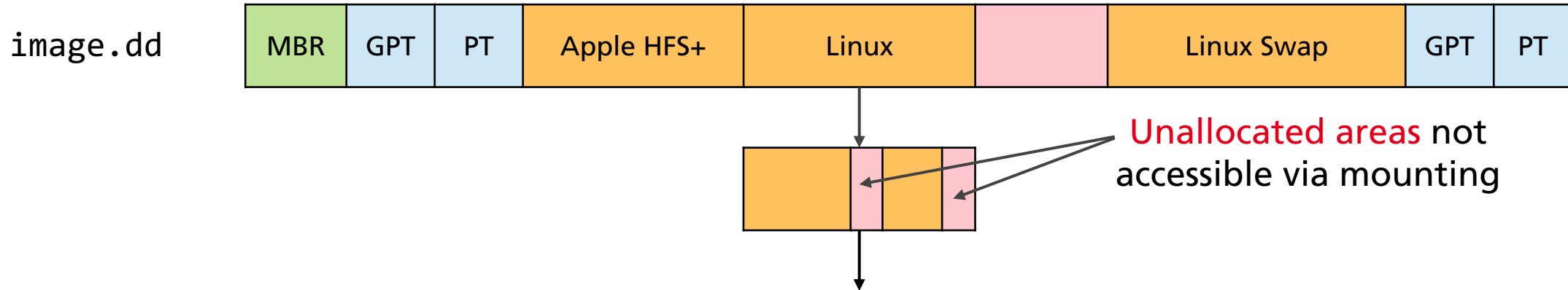
Analysis depends on the used implementation



Not all available information may be accessible



File System Analysis



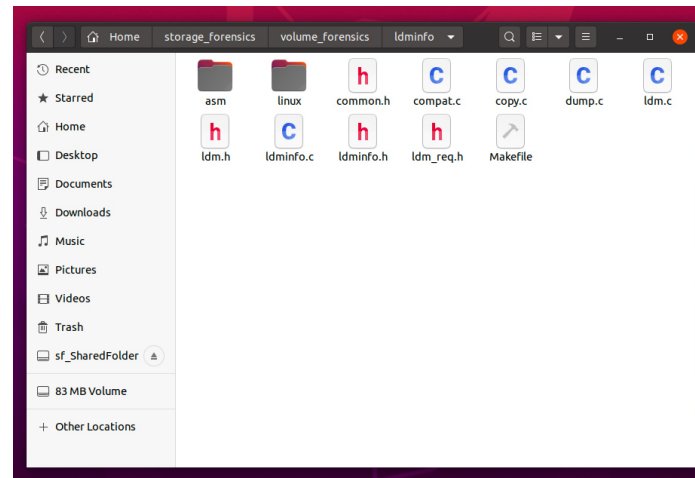
Analysis depends on the used implementation



Not all available information may be accessible



Deleted files are not mounted



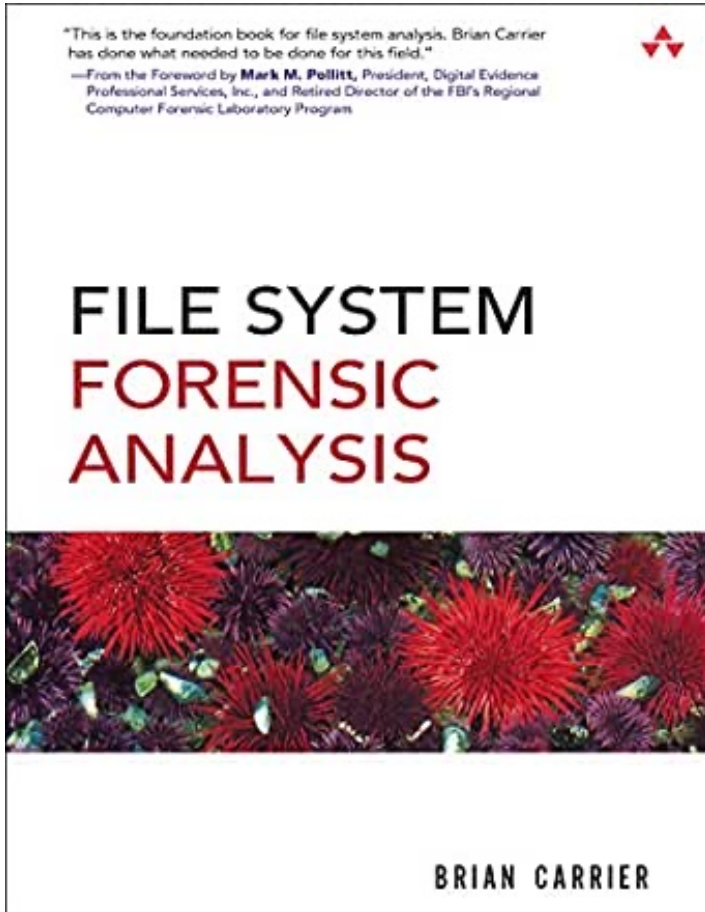
Unallocated areas not accessible via mounting

Use an independent analysis tool

Make sure all important information is easily accessible

Perform the analysis on the whole volume

The Sleuth Kit (TSK)



Open Source Digital Forensics



Autopsy® is an easy to use, GUI-based program that allows you to efficiently analyze hard drives and smart phones. It has a plug-in architecture that allows you to find add-on modules or develop custom modules in Java or Python.



The Sleuth Kit® is a collection of command line tools and a C library that allows you to analyze disk images and recover files from them. It is used behind the scenes in Autopsy and many other open source and commercial forensics tools.

Of course, TSK is not the only tool which performs file system analysis. But it's free, well-known and widely used!

These tools are used by thousands of users around the world and have community-based e-mail lists and forums. Commercial training, support, and custom development is available from Basis Technology.

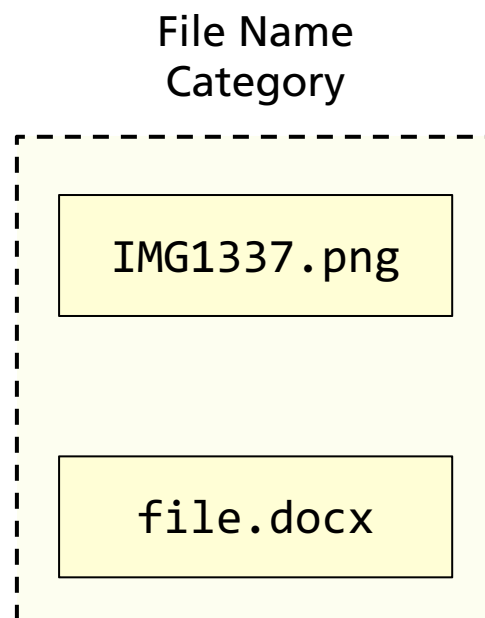
<https://www.sleuthkit.org>

File Systems

File Name Category

- Enables **users** to refer to a file using a name instead of an address
- Stores usually only a name and points to a **metadata entry**

1. How to **reference** a file?

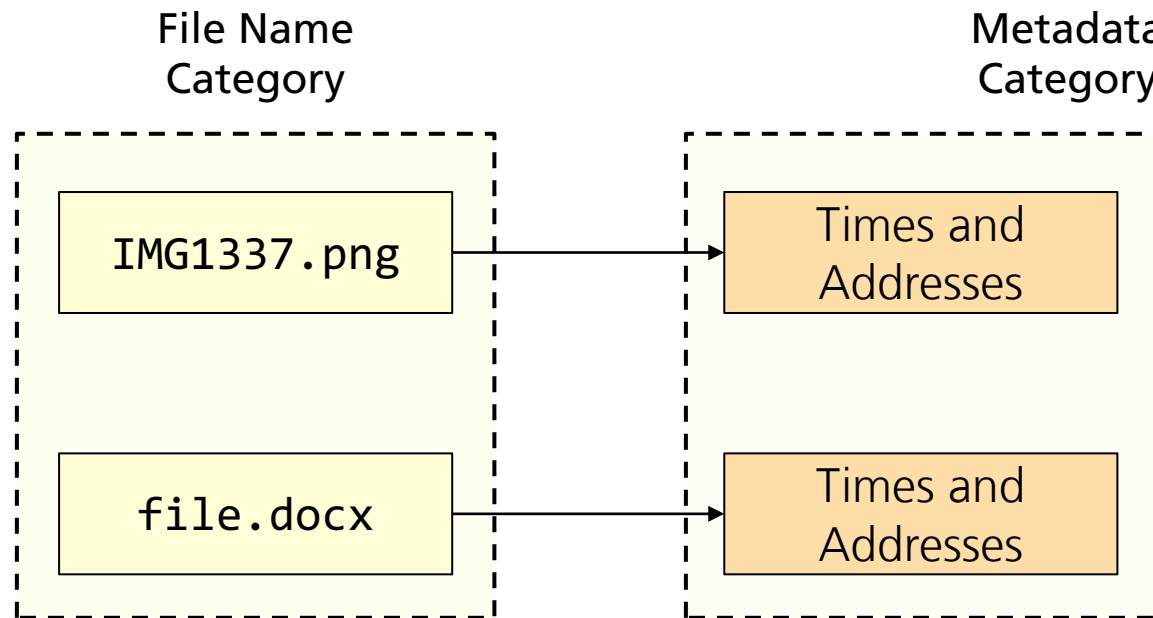


File Systems

Metadata Category

- Stores **metadata** about a file e.g. timestamps, sizes or attributes
- Often implemented by **metadata entries** with their own **allocation status**
- Often file name and metadata data are stored in **one place**
- Points to the **actual content** of a file

2. How to quickly find the **start** of a file and collect **metadata** (such as timestamps)?

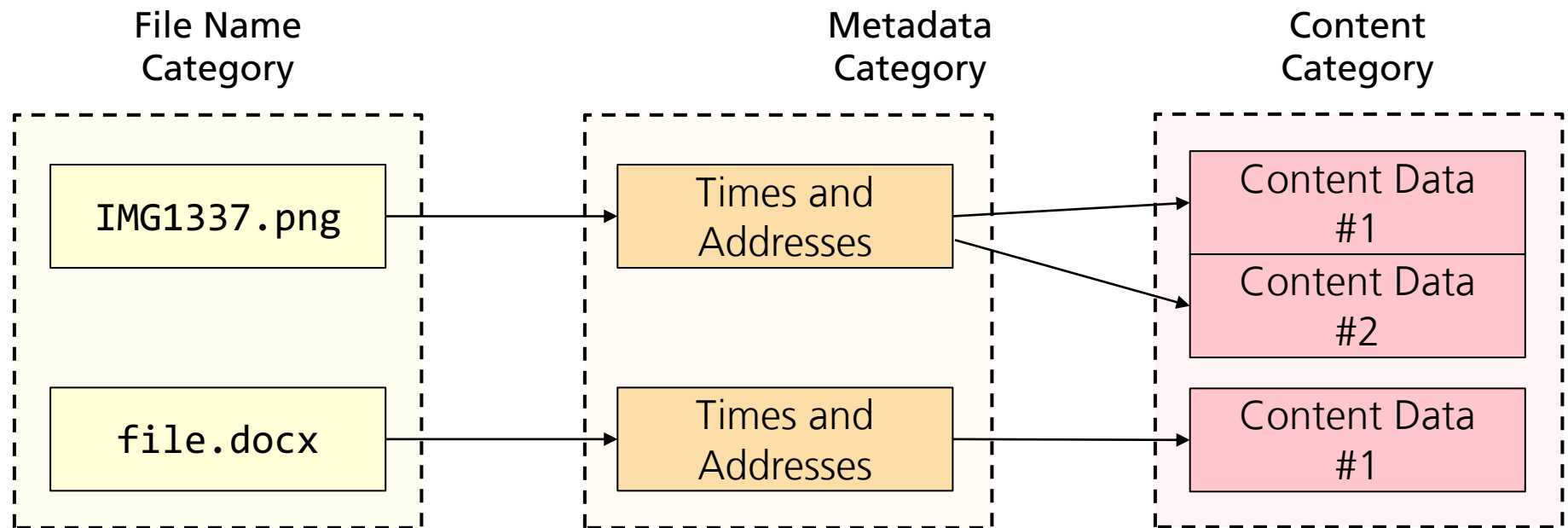


File Systems

Content Category

- Stores the **contents** of a **file** or **directory**
- Storage space is usually divided into **data units** (e.g. clusters or blocks) with their own **allocation status**
- It is sometimes not possible to store data units in **consecutive order**, which leads to **fragmentation**

3. How to organize **available space**?

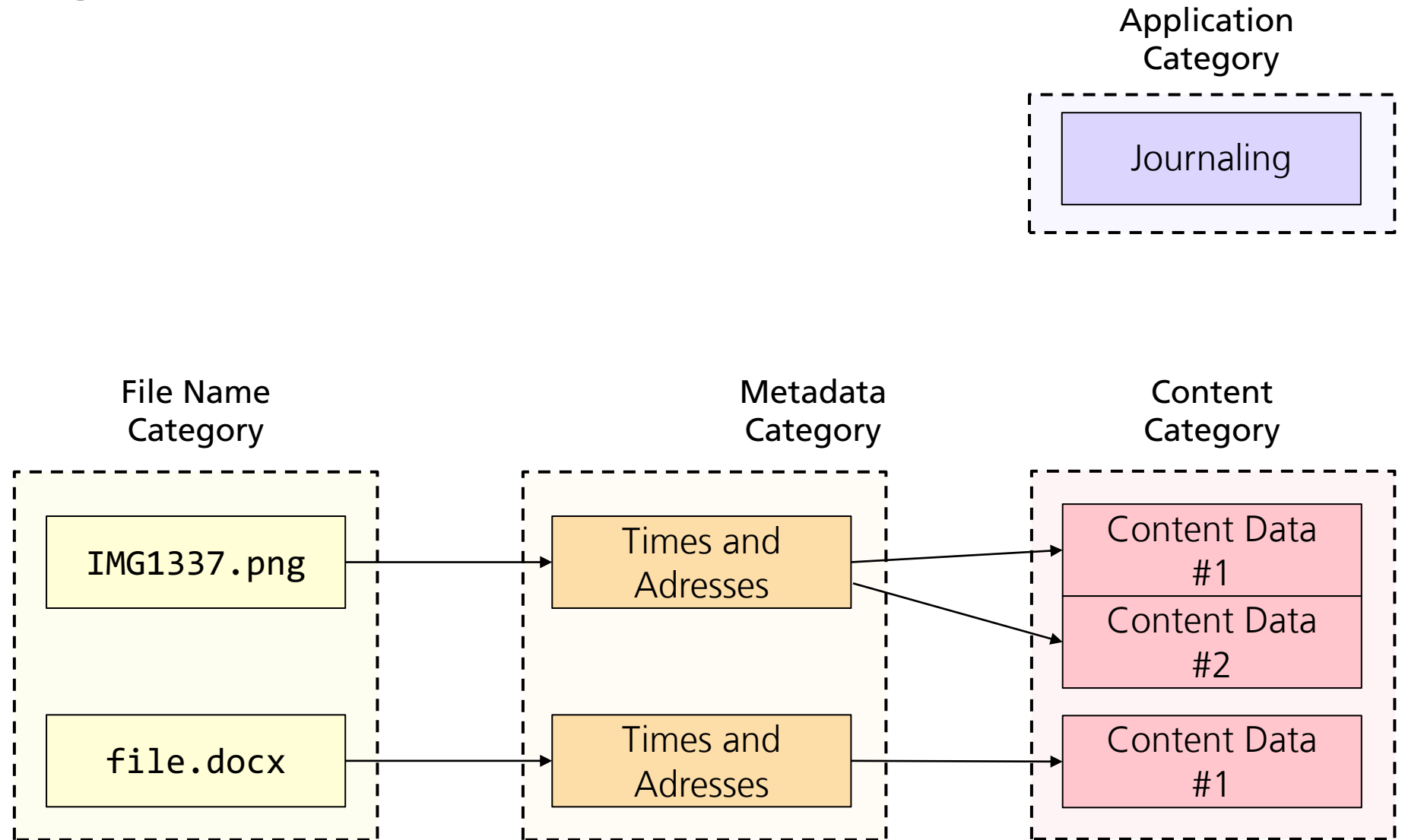


File Systems

Application Category

- Includes non-essential data used for special file system features, e.g. journaling or a quota

4. How to implement features such as logging?

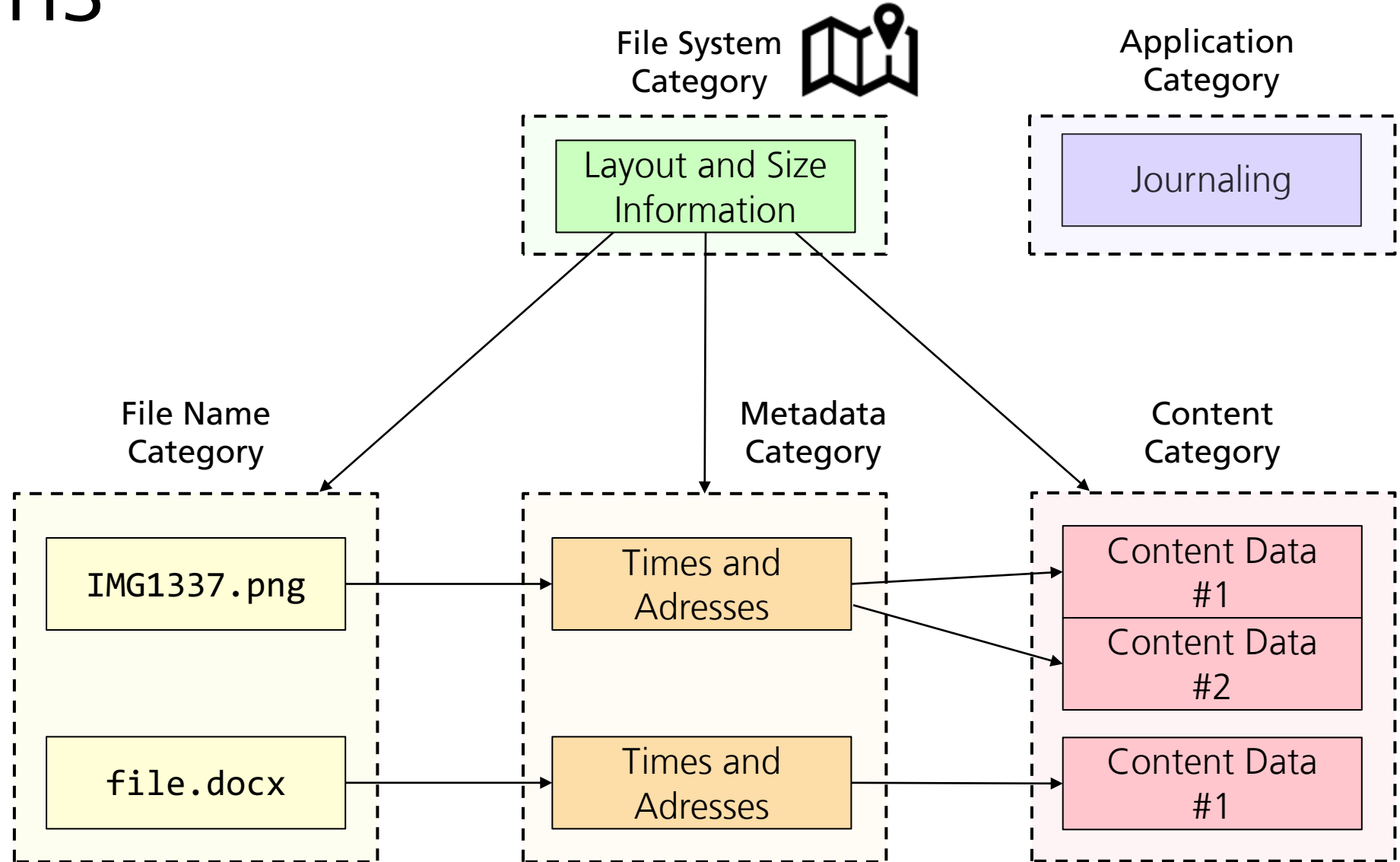


File Systems

File System Category

- Contains information about the **layout of a file system** including the location and sizes of its **important structures**
- Includes **metadata** about a file system (e.g. timestamps)
- **Crucial** data for **further analysis** of a file system

5. How to **organize** the aforementioned concepts?



Application Category
Journaling



File System Category

File Name Category

Metadata Category

Content Category

IMG1337.png

file.docx

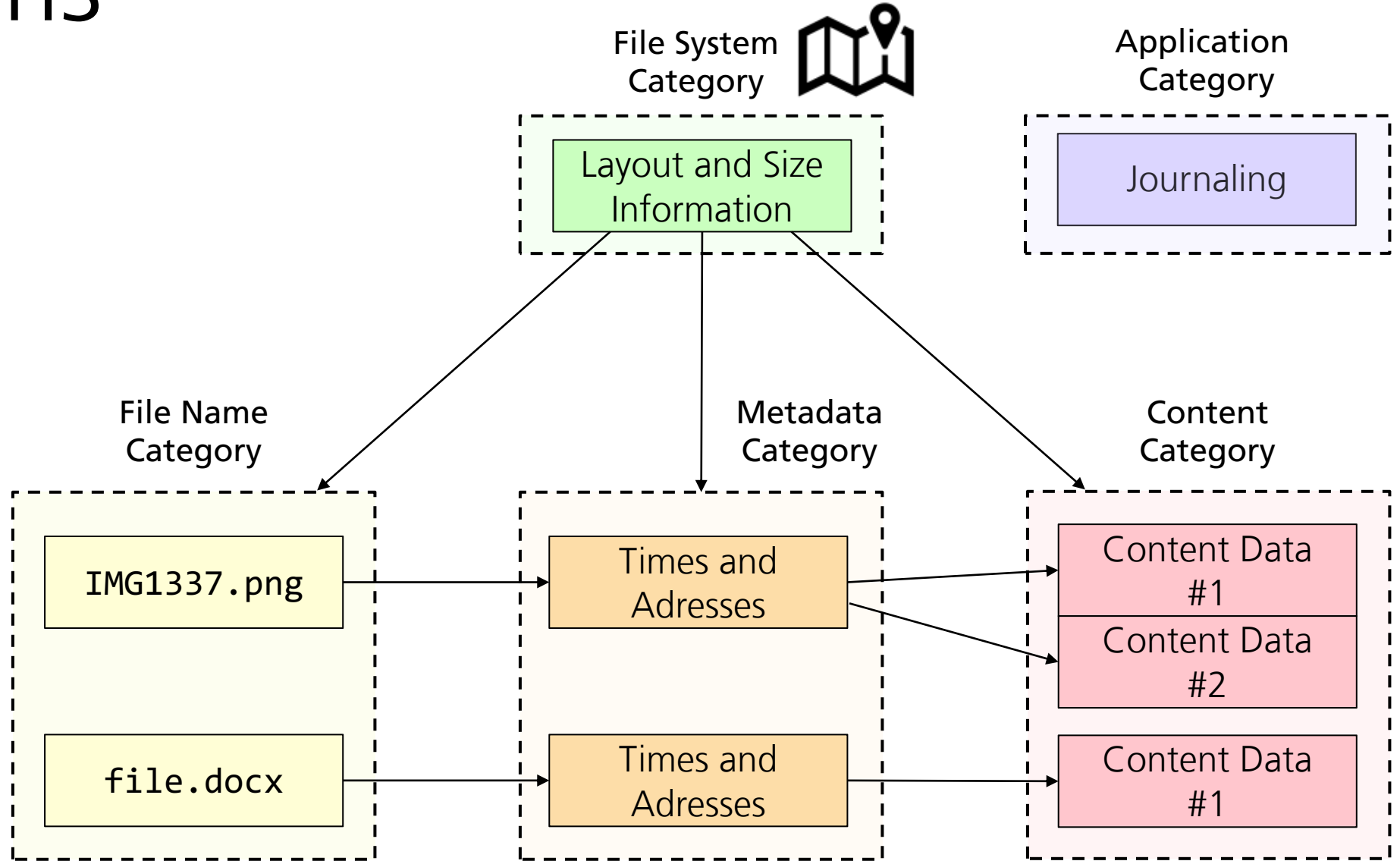
Times and Adresses

Times and Adresses

Content Data #1
Content Data #2

Content Data #1

File Systems



Dr. Brian Carrier

The Sleuth Kit

File System Tools

Fully Automated Tools

These tools integrate the volume and file system functionality. Instead of analyzing only a single file system, these tools take a disk image as input and identify the volumes and process the contents.

- [tsk_comparedir](#): Compares a local directory hierarchy with the contents of raw device (or disk image). This can be used to detect rootkits.
- [tsk_gettimes](#): Extracts all of the temporal data from the image to make a [timeline](#). Equivalent to running [fls](#) with the '-m' option.
- [tsk_loaddb](#): Loads the metadata from an image into a SQLite database. This allows other tools to be easily written in a variety of languages and give them access to the image contents.
- [tsk_recover](#): Extracts the unallocated (or allocated) files from a disk image to a local directory.

File System Layer Tools

These file system tools process general file system data, such as the layout, allocation structures, and boot blocks

- [fsstat](#): Shows file system details and statistics including layout, sizes, and labels.

File Name Layer Tools

These file system tools process the file name structures, which are typically located in the parent directory.

- [ffind](#): Finds allocated and unallocated file names that point to a given meta data structure.
- [fls](#): Lists allocated and deleted file names in a directory.

Meta Data Layer Tools

These file system tools process the meta data structures, which store the details about a file. Examples of this structure include directory entries in FAT, MFT entries in NTFS, and inodes in ExtX and UFS.

- [icat](#): Extracts the data units of a file, which is specified by its meta data address (instead of the file name).
- [ifind](#): Finds the meta data structure that has a given file name pointing to it or the meta data structure that points to a given data unit.
- [ils](#): Lists the meta data structures and their contents in a pipe delimited format.
- [istat](#): Displays the statistics and details about a given meta data structure in an easy to read format.

Data Unit Layer Tools

These file system tools process the [data units](#) where file content is stored. Examples of this layer include clusters in FAT and NTFS and blocks and fragments in ExtX and UFS.

- [blkcat](#): Extracts the contents of a given data unit.
- [blkls](#): Lists the details about data units and can extract the unallocated space of the file system.
- [blkstat](#): Displays the statistics about a given data unit in an easy to read format.
- [blkcalc](#): Calculates where data in the unallocated space image (from [blkls](#)) exists in the original image. This is used when evidence is found in unallocated space.

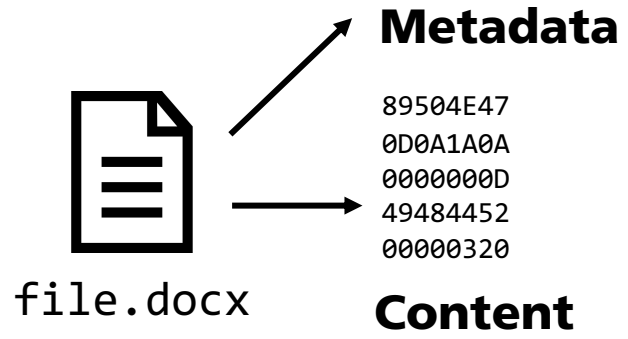
https://wiki.sleuthkit.org/index.php?title=TSK_Tool_Overview#File_System_Tools

File System Analysis

FAT File System

FAT file system

Name	Start	Size	Created on
file.docx	Data Unit 40	1400 Bytes	2015-07-10



Metadata

89504E47
0D0A1A0A
0000000D
49484452
00000320

Content

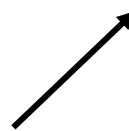
FAT file system

Name	Start	Size	Created on
file.docx	Data Unit 40	1400 Bytes	2015-07-10

Metadata in FAT is stored in "Directory Entry" structures



Metadata



file.docx

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00																
0x10																

Directory Entry

FAT file system

Name	Start	Size	Created on
file.docx	Data Unit 40	1400 Bytes	2015-07-10

Metadata



file.docx

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00	short name															
0x10																

Directory Entry

Up to 11 characters of the file's name in ASCII padded with spaces

FAT file system

Name	Start	Size	Created on
file.docx	Data Unit 40	1400 Bytes	2015-07-10

Metadata



file.docx

0xE5 This entry is free

0x00 This and following directory entries are free

0x05 The actual character is 0xE5

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00	short name															
0x10																

Directory Entry

FAT file system

Name	Start	Size	Created on
file.docx	Data Unit 40	1400 Bytes	2015-07-10

Metadata



file.docx

- 0x01 Read only
- 0x02 Hidden
- 0x04 Operating System file
- 0x08 Volume Label (should exist only once)
- 0x10 Directory
- 0x20 Archive

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00	short name											attr.				
0x10																

Directory Entry

File attributes

FAT file system

Name	Start	Size	Created on
file.docx	Data Unit 40	1400 Bytes	2015-07-10

Metadata



file.docx

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00	short name											attr.	res.	created time		
0x10	created date		date last accessed			time of last write			date of last write							

Directory Entry

FAT file system

Name	Start	Size	Created on
file.docx	Data Unit 40	1400 Bytes	2015-07-10

Metadata



file.docx

byte offset	0x00	0x01	0x02	0x03
0x00				
0x10	created date		date last accessed	

Directory Entry

Date Format. A FAT directory entry date stamp is a 16-bit field that is basically a date relative to the MS-DOS epoch of 01/01/1980. Here is the format (bit 0 is the LSB of the 16-bit word, bit 15 is the MSB of the 16-bit word):

Bits 0–4: Day of month, valid value range 1-31 inclusive.

Bits 5–8: Month of year, 1 = January, valid value range 1–12 inclusive.

Bits 9–15: Count of years from 1980, valid value range 0–127 inclusive (1980–2107).

Time Format. A FAT directory entry time stamp is a 16-bit field that has a granularity of 2 seconds. Here is the format (bit 0 is the LSB of the 16-bit word, bit 15 is the MSB of the 16-bit word).

Bits 0–4: 2-second count, valid value range 0–29 inclusive (0 – 58 seconds).

Bits 5–10: Minutes, valid value range 0–59 inclusive.

Bits 11–15: Hours, valid value range 0–23 inclusive.

The valid time range is from Midnight 00:00:00 to 23:59:58.

FAT file system

Analysis: Timestamps



- **Timestamps** provide **valuable information** during an investigation
- **However:** It is always important to know, what **they mean** i.e. when exactly they are updated
- **Example:** Findings for the **Creation Time** in **FAT** by Brian Carrier:
 - *„The created time is set when Windows allocates a new directory entry for a new file.“*
 - *„[...] if the OS allocates a new directory entry for an existing file, even if original location was on a different disk, the original creation time is kept.“*
 - *„There is one known exception to this rule, if the move is done from the command line of a 2000/XP system to a different volume.“*
- **Summary:** **Time Value Updating** depends on a **combination** of file system and operating system (and tools used) and has to be evaluated in order to know the exact meaning of available time stamps

FAT file system

Analysis: Timestamps

- **Timestamps** provide **valuable information** during an
- **However:** It is always important to know, what **they** they are updated
- **Example:** Findings for the **Creation Time** in **FAT** by B
 - *„The created time is set when Windows allocates new file.“*
 - *„[...] if the OS allocates a new directory entry for original location was on a different disk, the origi*
 - *„There is one known exception to this rule, if the command line of a 2000/XP system to a different*
- **Summary:** **Time Value Updating** depends on a **comb** operating system (and tools used) and has to be eval exact meaning of available time stamps



SANS Institute Information Security Reading Room

Filesystem Timestamps: What Makes Them Tick?

Tony Knutson

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FAT file system

Name	Start	Size	Created on
file.docx	Data Unit 40	1400 Bytes	2015-07-10

Metadata



file.docx

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00	short name											attr.	res.	created time		
0x10	created date		date last accessed			time of last write			date of last write							

Directory Entry

FAT file system

Name	Start	Size	Created on
file.docx	Data Unit 40	1400 Bytes	2015-07-10

Metadata



file.docx

Number of first cluster
(high bytes are 0 for FAT12/16)

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00	short name										attr.	res.	created time			
0x10	created date	date last accessed		first cluster (high bytes)		time of last write		date of last write		first cluster (low bytes)						

Directory Entry

FAT file system

Name	Start	Size	Created on
file.docx	Data Unit 40	1400 Bytes	2015-07-10

Metadata



file.docx

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00	short name											attr.	res.	created time		
0x10	created date	date last accessed		first cluster (high bytes)		time of last write		date of last write		first cluster (low bytes)		file size				

Directory Entry

File size in bytes (max. 4 GB!)

FAT file system

Name	Start	Size	Created on
file-with-a-long-name.docx	Data Unit 40	1400 Bytes	2015-07-10

Metadata

What about „really long“ file names?



file*.docx

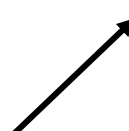
byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00	short name											attr.	res.	created time		
0x10	created date	date last accessed		first cluster (high bytes)		time of last write		date of last write		first cluster (low bytes)		file size				

Directory Entry

FAT file system

Name	Start	Size	Created on
file-with-a-long-name.docx	Data Unit 40	1400 Bytes	2015-07-10

Metadata



file*.docx

For longer names, the file name is split up and stored in multiple Long File Name (LFN) directory entries (in addition to a normal directory entry)!

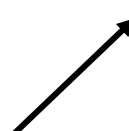
byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00	characters 1 – 5															
0x10	characters 6 – 11											characters 12 – 13				

LFN Directory Entry

FAT file system

Name	Start	Size	Created on
file-with-a-long-name.docx	Data Unit 40	1400 Bytes	2015-07-10

Metadata



file*.docx

Used to order multiple LFN entries.
Last part is masked with 0x40 and should be the first LFN entry.

Sequence #	Entry data
0x42	ong-name.docx
0x01	file-with-a-l
-	Default Directory Entry

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00	seq. #	characters 1 – 5														
0x10	characters 6 – 11										characters 12 – 13					

LFN Directory Entry

FAT file system

Name	Start	Size	Created on
file-with-a-long-name.docx	Data Unit 40	1400 Bytes	2015-07-10

Metadata



file*.docx

Zero indicates an entry that is part of a long file name

Special attribute value

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00	seq. #	characters 1 – 5										0x0F	0x00	check-sum		
0x10	characters 6 – 11												characters 12 – 13			

LFN Directory Entry

FAT file system

Name	Start	Size	Created on
file-with-a-long-name.docx	Data Unit 40	1400 Bytes	2015-07-10

Metadata



file*.docx

Checksum of the corresponding
short file name

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00	seq. #	characters 1 – 5										0x0F	0x00	check- sum		
0x10	characters 6 – 11										characters 12 – 13					

LFN Directory Entry

FAT file system

Name	Start	Size	Created on
file-with-a-long-name.docx	Data Unit 40	1400 Bytes	2015-07-10

Metadata



file*.docx

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00	seq. #	characters 1 – 5										0x0F	0x00	check- sum		
0x10	characters 6 – 11										0x0000		characters 12 – 13			

LFN Directory Entry

FAT file system

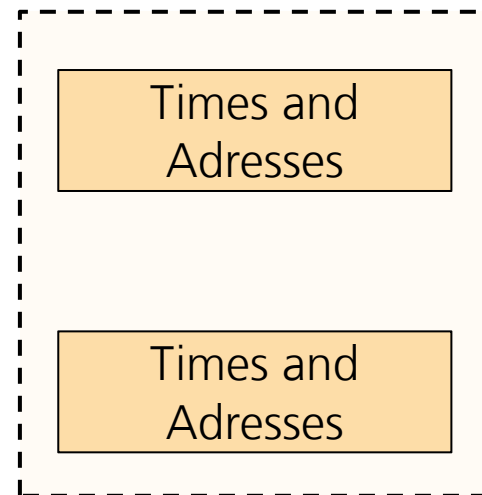
Name	Start	Size	Created on
file.docx	Data Unit 40	1400 Bytes	2015-07-10

Metadata



file.docx

Metadata
Category



Directory entries belong
to the metadata category...

FAT file system

Name	Start	Size	Created on
file.docx	Data Unit 40	1400 Bytes	2015-07-10

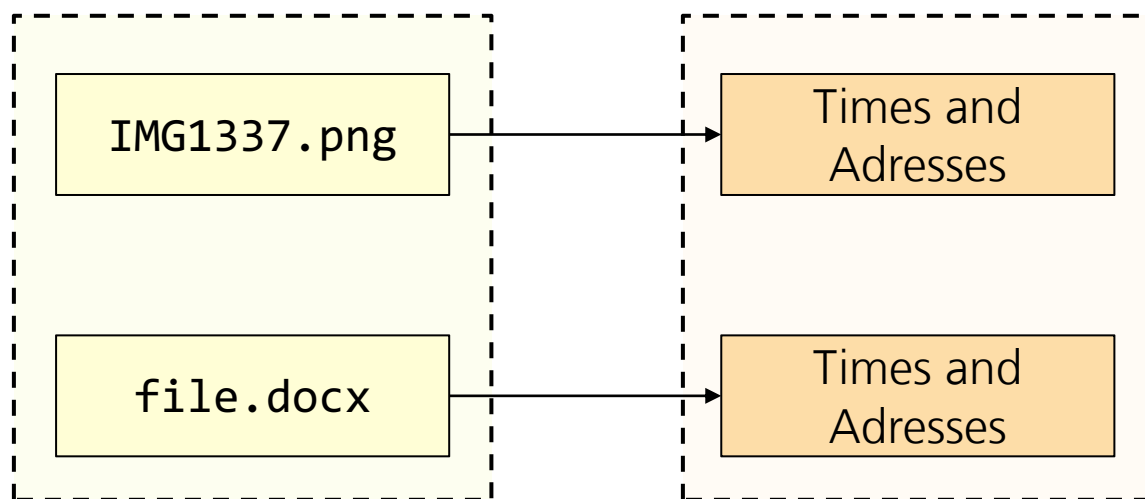
Metadata



file.docx

File Name Category

Metadata Category



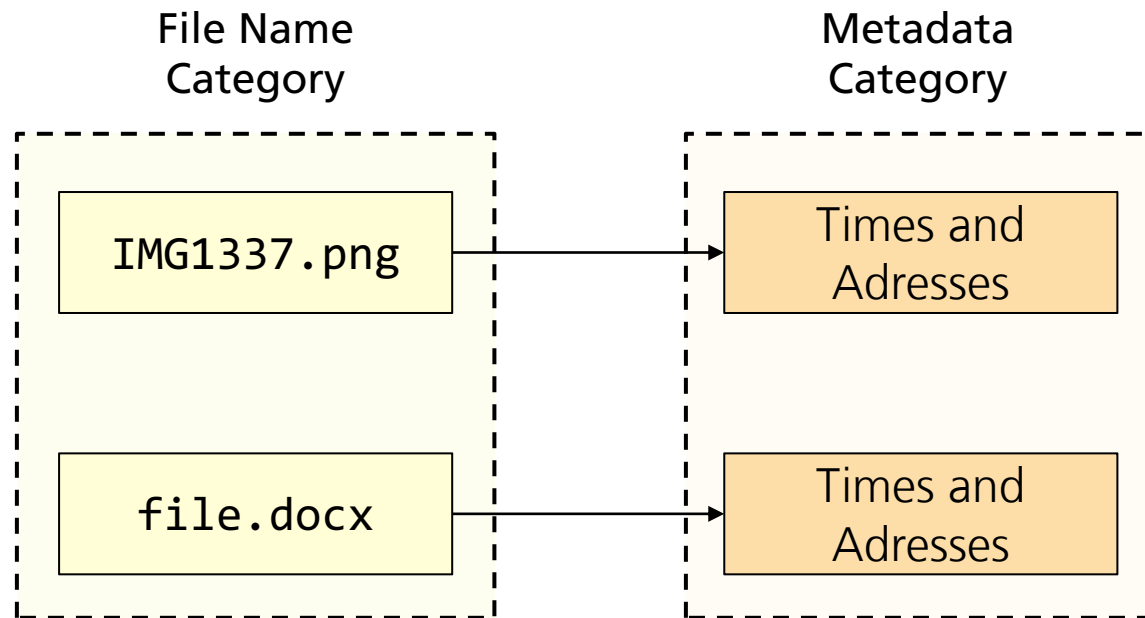
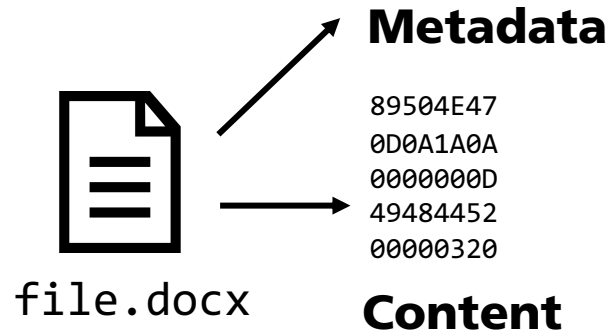
...but also partially contain file name data.

LFNs are pure file name data

Directory entries belong to the metadata category...

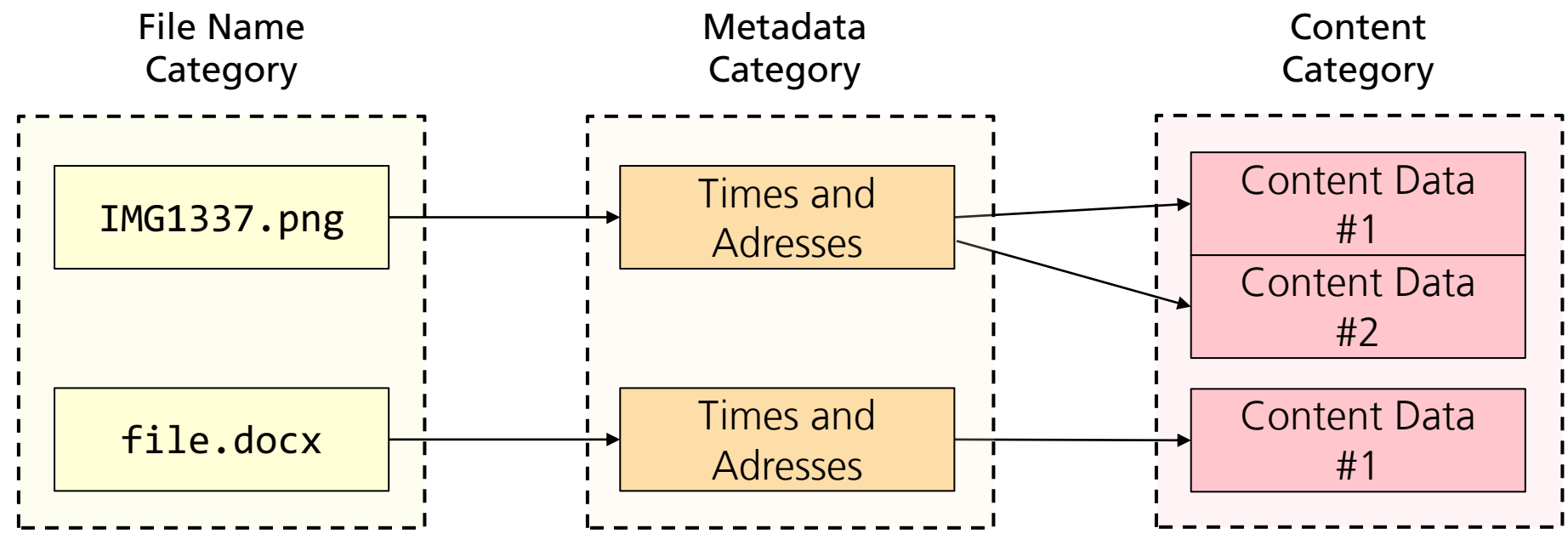
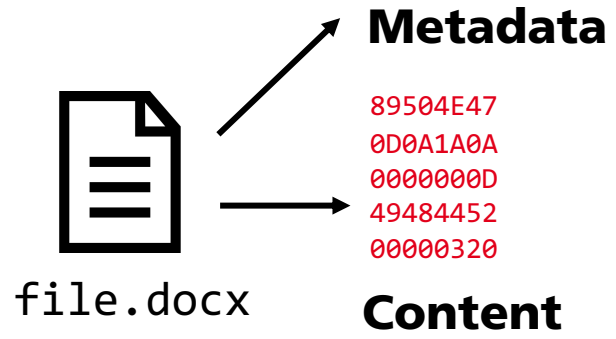
FAT file system

Name	Start	Size	Created on
file.docx	Data Unit 40	1400 Bytes	2015-07-10



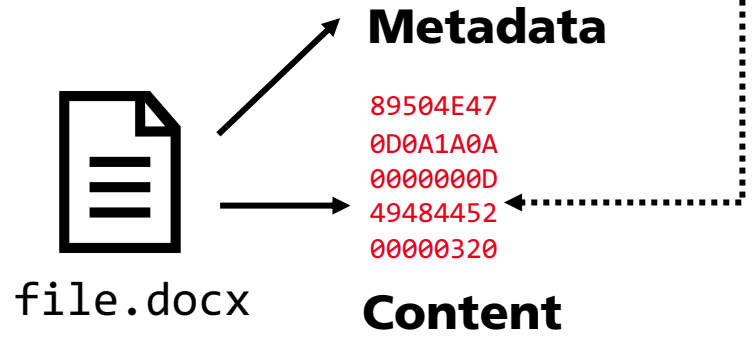
FAT file system

Name	Start	Size	Created on
file.docx	Data Unit 40	1400 Bytes	2015-07-10



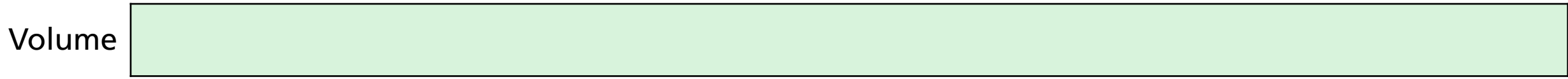
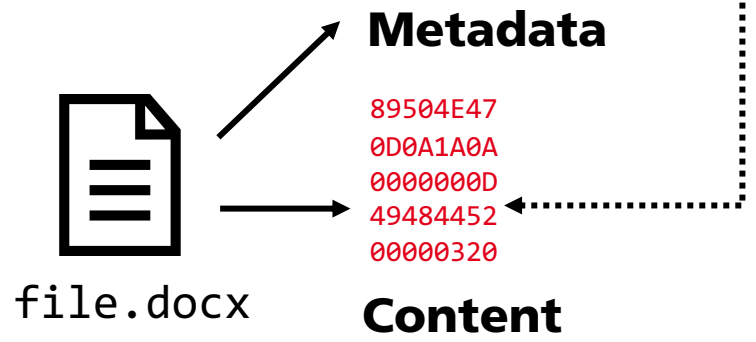
FAT file system

Name	Start	Size	Created on
file.docx	Data Unit 40	1400 Bytes	2015-07-10



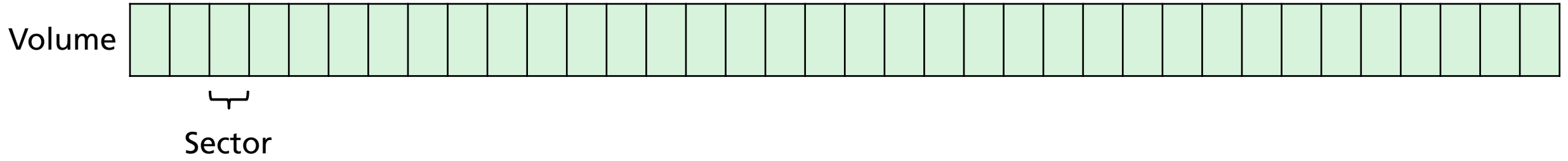
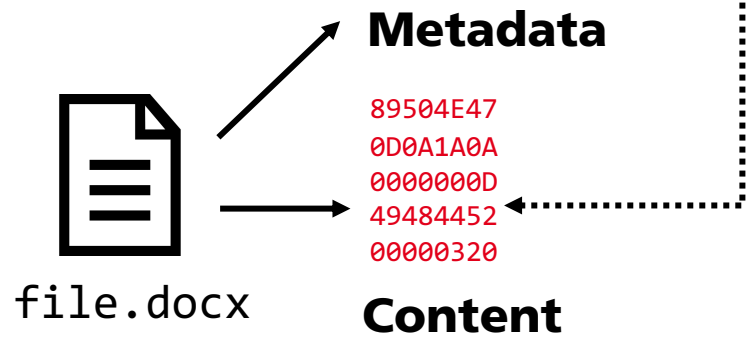
FAT file system

Name	Start	Size	Created on
file.docx	Data Unit 40	1400 Bytes	2015-07-10



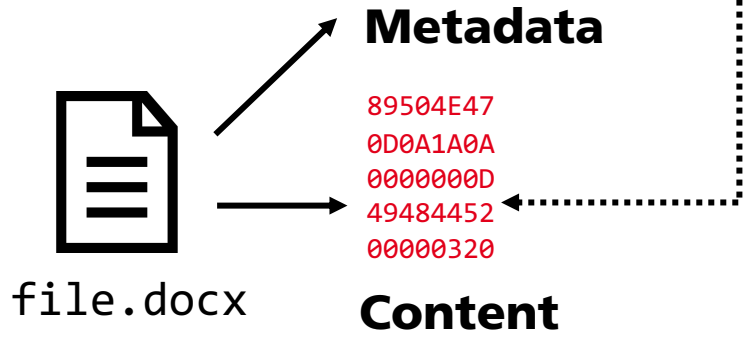
FAT file system

Name	Start	Size	Created on
file.docx	Data Unit 40	1400 Bytes	2015-07-10

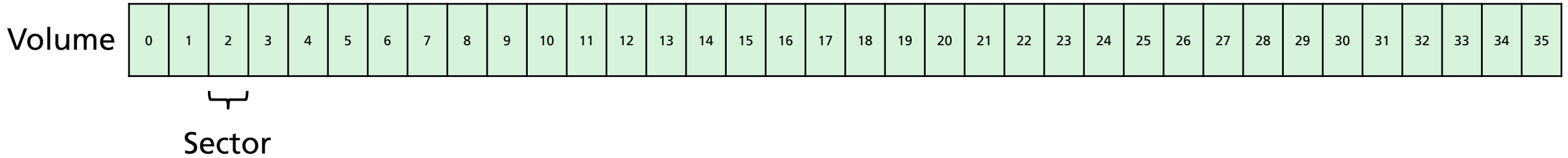


FAT file system

Name	Start	Size	Created on
file.docx	Data Unit 40	1400 Bytes	2015-07-10

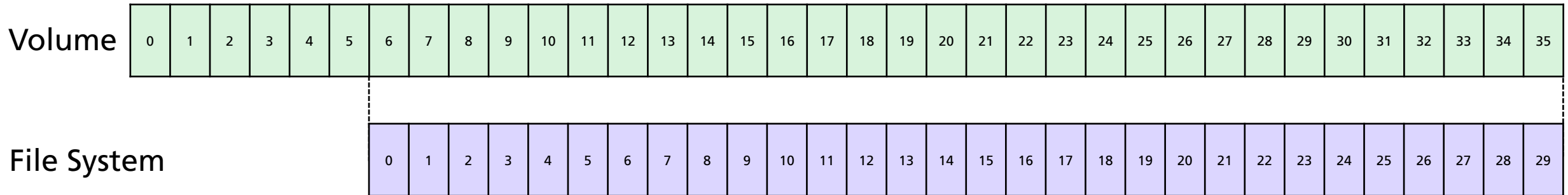
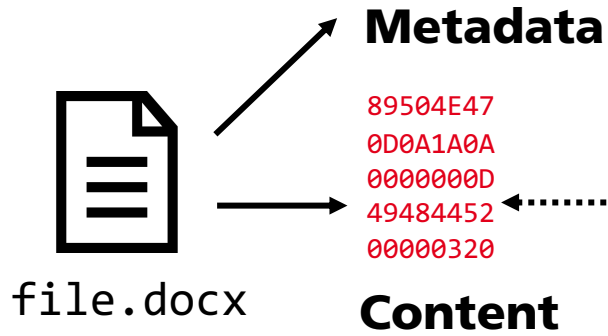


Sectors can be addresses by their "Logical Volume Addresses"



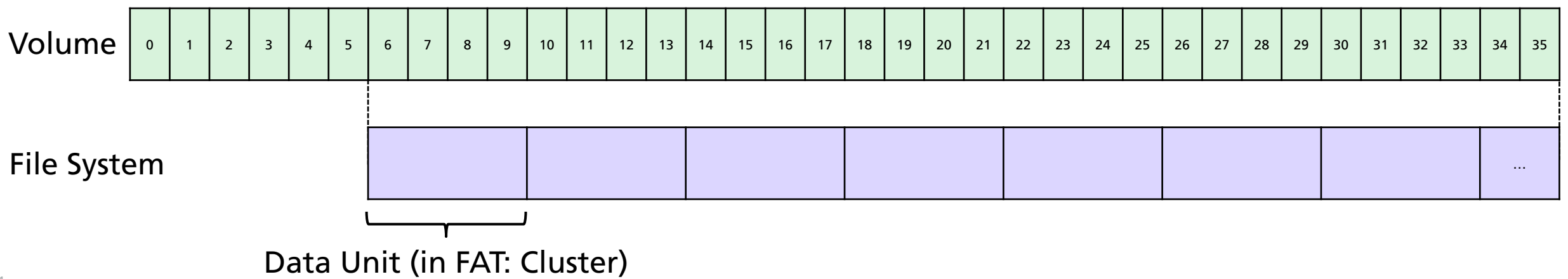
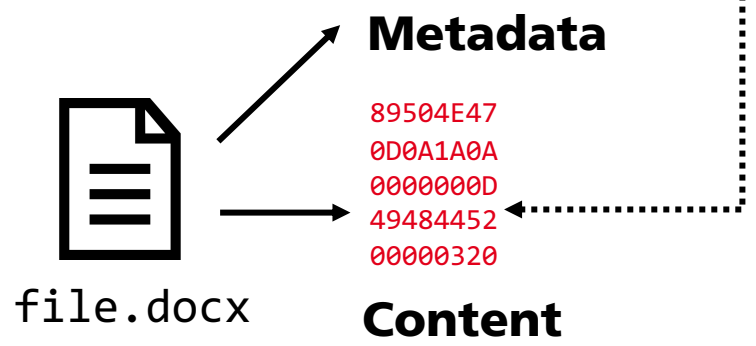
FAT file system

Name	Start	Size	Created on
file.docx	Data Unit 40	1400 Bytes	2015-07-10



FAT file system

Name	Start	Size	Created on
file.docx	Data Unit 40	1400 Bytes	2015-07-10



FAT file system

Name	Start	Size	Created on
file.docx	Data Unit 40	1400 Bytes	2015-07-10



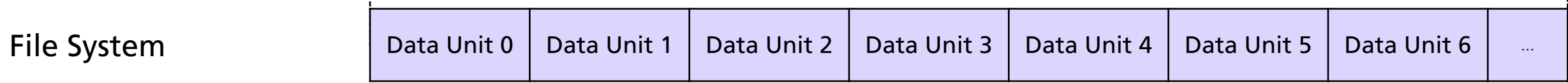
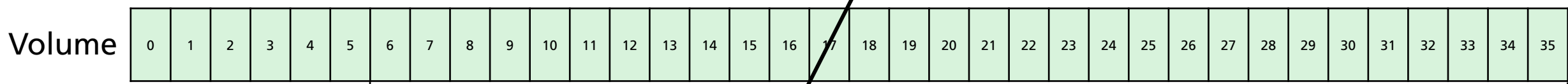
file.docx

Metadata

89504E47
 0D0A1A0A
 0000000D
 49484452
 00000320

Content

Data Units can be addresses by their "Logical File System Addresses"



Data Unit (in FAT: Cluster)

FAT file system

mkfs.fat

mkfs.fat(8) - Linux man page

Name

mkfs.fat - create an MS-DOS FAT filesystem

Synopsis

mkfs.fat [OPTIONS] DEVICE [BLOCK-COUNT]

Description

mkfs.fat is used to create a FAT filesystem on a device or in an image file.

[...]

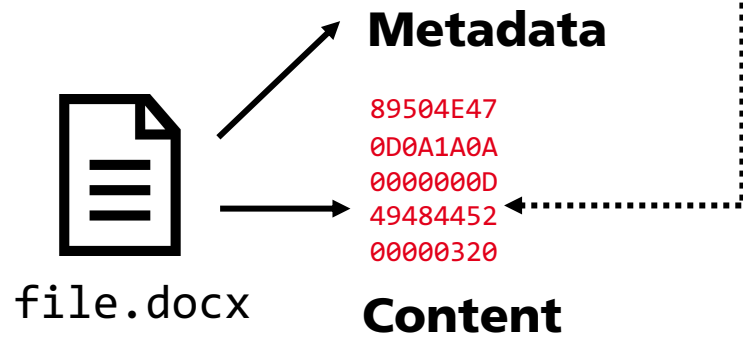
Options

-s *SECTORS-PER-CLUSTER* Specify the number of **disk sectors per cluster**. Must be a power of 2, i.e. 1, 2, 4, 8, ... 128.

-S *LOGICAL-SECTOR-SIZE* Specify the **number of bytes per logical sector**. Must be a power of 2 and greater than or equal to 512, i.e. 512, 1024, 2048, 4096, 8192, 16384, or 32768. Values larger than 4096 are not conforming to the FAT file system specification and may not work everywhere.

FAT file system

Name	Start	Size	Created on
file.docx	Data Unit 40	1400 Bytes	2015-07-10



Used to store data for files
and directories

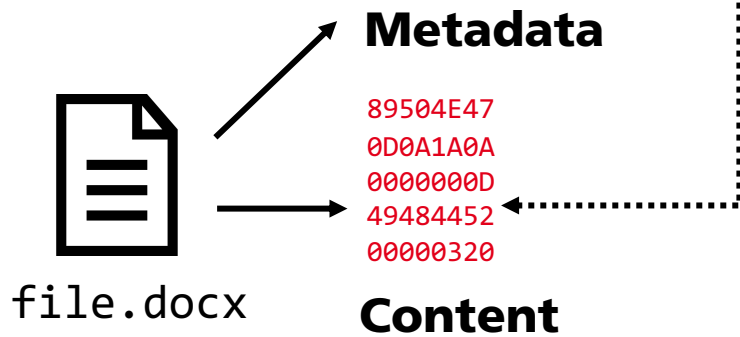
File System



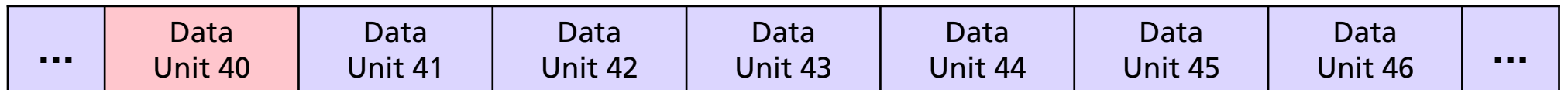
Data Area

FAT file system

Name	Start	Size	Created on
file.docx	Data Unit 40	1400 Bytes	2015-07-10

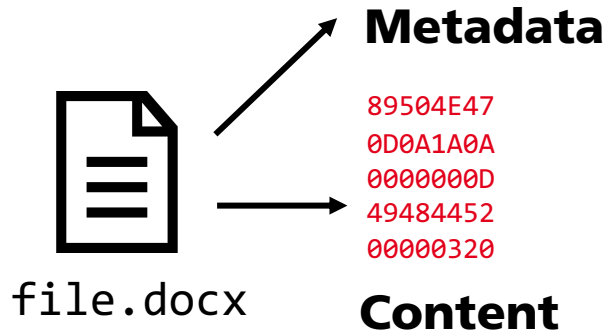


File System



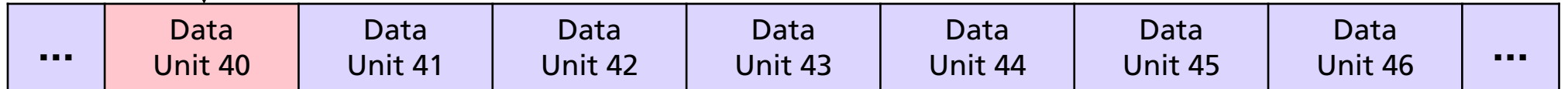
FAT file system

Name	Start	Size	Created on
file.docx	Data Unit 40	1400 Bytes	2015-07-10



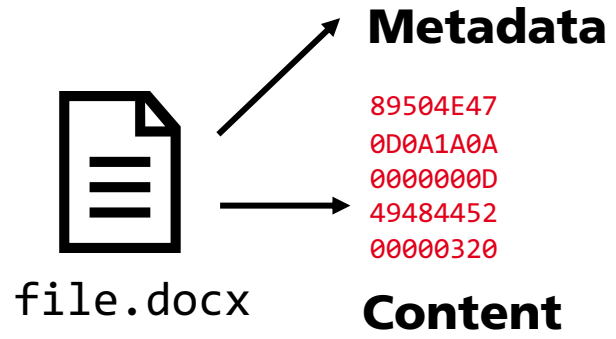
What if the file is smaller than a data unit?
=> Doesn't matter!

File System



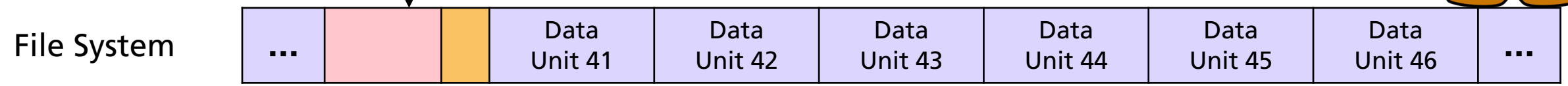
FAT file system

Name	Start	Size	Created on
file.docx	Data Unit 40	1400 Bytes	2015-07-10



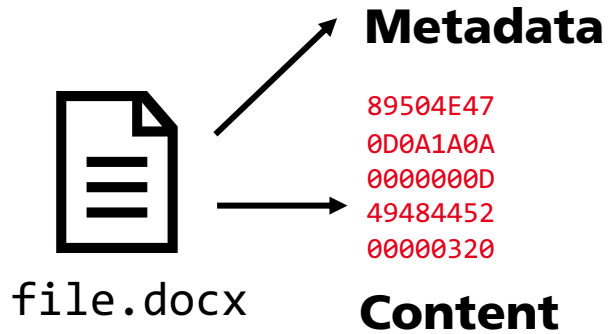
What if the file is smaller than a data unit?
=> Doesn't matter!

Results in **file slack** between the end of the file and the next data unit!



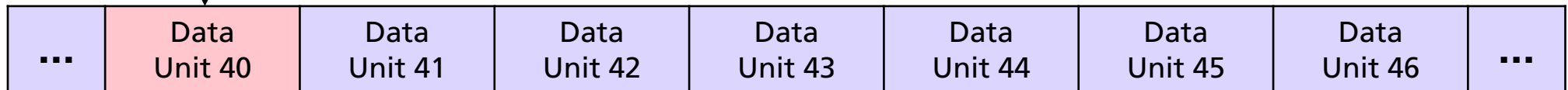
FAT file system

Name	Start	Size	Created on
file.docx	Data Unit 40	1400 Bytes	2015-07-10



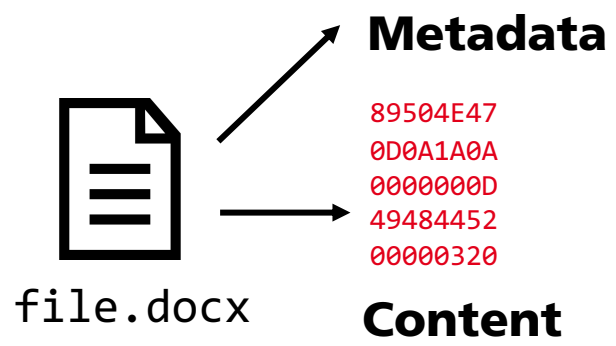
What about files larger than a data unit?

File System



FAT file system

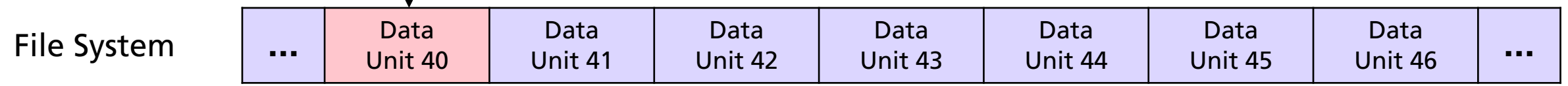
Name	Start	Size	Created on
file.docx	Data Unit 40	1400 Bytes	2015-07-10



The File Allocation Table is used to enable **cluster chains/runs**.

Entry	Value
...	...
40	41
41	45
42	43
43	EOF
44	0
45	EOF
...	...

FAT



FAT file system

Name	Start	Size	Created on
file.docx	Data Unit 40	1400 Bytes	2015-07-10

Entry	Value
...	...
40	41
41	45
42	43
43	EOF
44	0
45	EOF
...	...

Metadata

89504E47
 0D0A1A0A
 0000000D
 49484452
 00000320

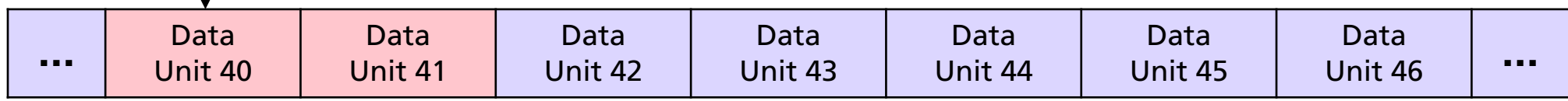
Content

The File Allocation Table is used to enable cluster chains/runs.

FAT

file.docx

File System



FAT file system

Name	Start	Size	Created on
file.docx	Data Unit 40	1400 Bytes	2015-07-10

Entry	Value
...	...
40	41
41	45
42	43
43	EOF
44	0
45	EOF
...	...



file.docx

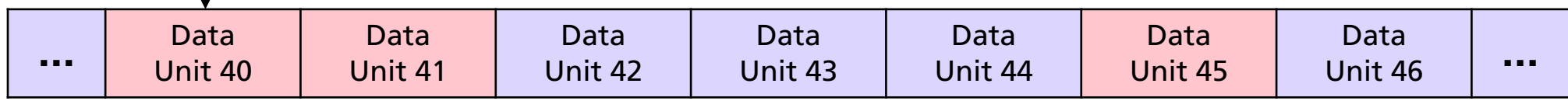
Metadata

89504E47
 0D0A1A0A
 0000000D
 49484452
 00000320

Content

The File Allocation Table is used to enable **cluster chains/runs**.

File System



FAT

FAT file system

Name	Start	Size	Created on
file.docx	Data Unit 40	1400 Bytes	2015-07-10

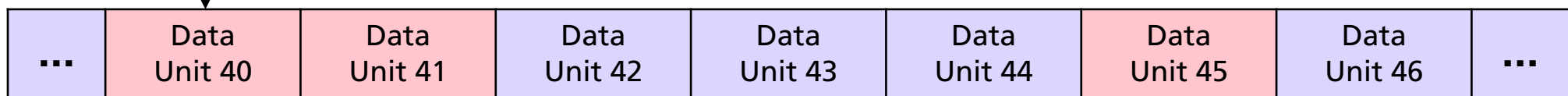
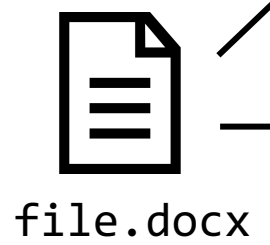
Entry	Value
...	...
40	41
41	45
42	43
43	EOF
44	0
45	EOF
...	...

Metadata

89504E47
 0D0A1A0A
 0000000D
 49484452
 00000320

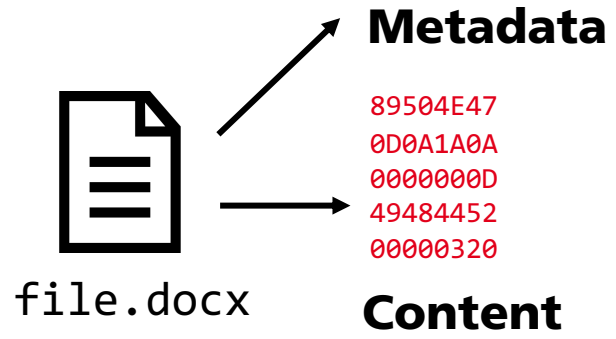
Content

The File Allocation Table is used to enable cluster chains/runs.



FAT file system

Name	Start	Size	Created on
file.docx	Data Unit 40	1400 Bytes	2015-07-10

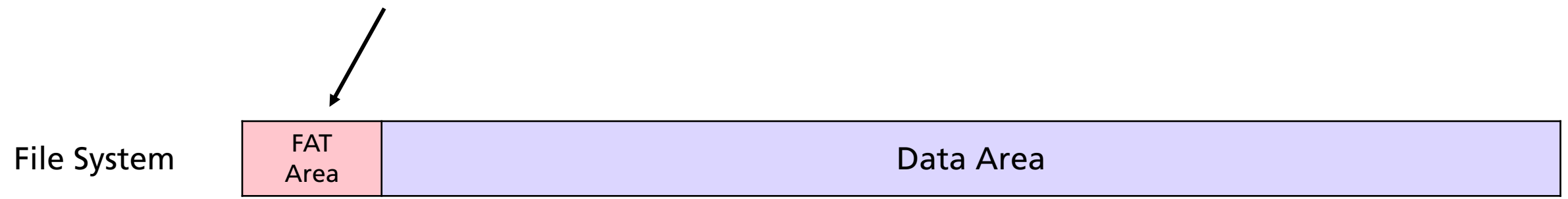


The File Allocation Table is used to enable **cluster chains/runs**.

Entry	Value
...	...
40	41
41	45
42	43
43	EOF
44	0
45	EOF
...	...

FAT

Stores the File Allocation Table(s)





FAT file system

FAT (FAT32)

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00																
0x10																
...																



FAT file system

FAT (FAT32)

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f	
0x00	reserved				reserved												
0x10																	
...																	

FAT file system

FAT (FAT32)

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f	
0x00	0x0FFFFFF8				0x0FFFFFFF												
0x10																	
...																	

Some special values, we don't care about!



FAT file system

FAT (FAT32)

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00	0x0FFFFFF8				0x0FFFFFFF				Cluster following cluster 2				Cluster following cluster 3			
0x10	Cluster following cluster 4				Cluster following cluster 5				Cluster following cluster 6				Cluster following cluster 7			
...																

0x?0000000

Free cluster marker



Only 28 of the available 32 bits are used!

FAT file system

FAT (FAT32)

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00	0x0FFFFFF8				0x0FFFFFFF				Cluster following cluster 2				Cluster following cluster 3			
0x10	Cluster following cluster 4				Cluster following cluster 5				Cluster following cluster 6				Cluster following cluster 7			
...																

0x?0000000

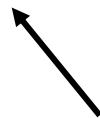
Free cluster marker

0x?FFFFFF7

Bad cluster marker

0x?FFFFFF8 or higher

End of file marker



Set to 0x0FFFFFFF by Microsoft FAT drivers



FAT file system

Example: FAT (FAT32)

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00	0x0FFFFFFF8				0x0FFFFFFF				0x00000005				0x00000004			
0x10	0x00000007				0x0FFFFFFF8				0x0000000A				0x00000006			
0x20	0x00000000				0x0000F541				0x0FFFFFFF8				0x00000000			

IMG9109.jpeg
start cluster: 3

file.bin
start cluster: 2

FAT file system

Example: FAT (FAT32)

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00	0x0FFFFFF8				0x0FFFFFFF				0x00000005				0x00000004			
0x10	0x00000007				0x0FFFFFF8				0x0000000A				0x00000006			
0x20	0x00000000				0x0000F541				0x0FFFFFF8				0x00000000			

IMG9109.jpeg
start cluster: 3

file.bin
start cluster: 2

FAT file system

Example: FAT (FAT32)

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00	0x0FFFFFF8				0x0FFFFFFF				0x00000005				0x00000004			
0x10	0x00000007				0x0FFFFFF8				0x0000000A				0x00000006			
0x20	0x00000000				0x0000F541				0x0FFFFFF8				0x00000000			

IMG9109.jpeg
start cluster: 3

3

file.bin
start cluster: 2

FAT file system

Example: FAT (FAT32)

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00	0x0FFFFFF8				0x0FFFFFFF				0x00000005				0x00000004			
0x10	0x00000007				0x0FFFFFF8				0x0000000A				0x00000006			
0x20	0x00000000				0x0000F541				0x0FFFFFF8				0x00000000			

IMG9109.jpeg
start cluster: 3

3 → 4

file.bin
start cluster: 2

FAT file system

Example: FAT (FAT32)

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00	0x0FFFFFF8				0x0FFFFFFF				0x00000005				0x00000004			
0x10	0x00000007				0x0FFFFFF8				0x0000000A				0x00000006			
0x20	0x00000000				0x0000F541				0x0FFFFFF8				0x00000000			

IMG9109.jpeg
start cluster: 3

3 → 4 → 7

file.bin
start cluster: 2

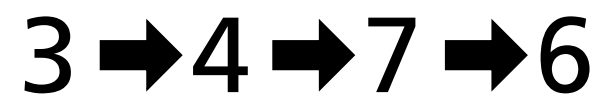


FAT file system

Example: FAT (FAT32)

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00	0x0FFFFFF8				0x0FFFFFFF				0x00000005				0x00000004			
0x10	0x00000007				0x0FFFFFF8				0x0000000A				0x00000006			
0x20	0x00000000				0x0000F541				0x0FFFFFF8				0x00000000			

IMG9109.jpeg
start cluster: 3



file.bin
start cluster: 2

FAT file system

Example: FAT (FAT32)

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00	0x0FFFFFF8				0x0FFFFFFF				0x00000005				0x00000004			
0x10	0x00000007				0x0FFFFFF8				0x0000000A				0x00000006			
0x20	0x00000000				0x0000F541				0x0FFFFFF8				0x00000000			

IMG9109.jpeg
start cluster: 3

3 → 4 → 7 → 6 → 10

file.bin
start cluster: 2

FAT file system

Example: FAT (FAT32)

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00	0x0FFFFFF8				0x0FFFFFFF				0x00000005				0x00000004			
0x10	0x00000007				0x0FFFFFF8				0x0000000A				0x00000006			
0x20	0x00000000				0x0000F541				0x0FFFFFF8				0x00000000			

IMG9109.jpeg
start cluster: 3

3 → 4 → 7 → 6 → 10

file.bin
start cluster: 2

2



FAT file system

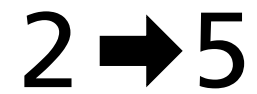
Example: FAT (FAT32)

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00	0x0FFFFFF8				0x0FFFFFFF				0x00000005				0x00000004			
0x10	0x00000007				0x0FFFFFF8				0x0000000A				0x00000006			
0x20	0x00000000				0x0000F541				0x0FFFFFF8				0x00000000			

IMG9109.jpeg
start cluster: 3



file.bin
start cluster: 2





FAT file system

Analysis: Unallocated Area

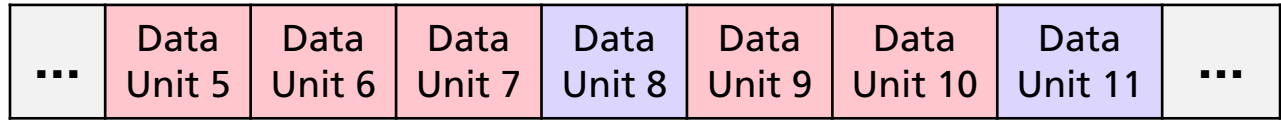


byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00	0x0FFFFFFF8				0x0FFFFFFF				0x00000005				0x00000004			
0x10	0x00000007				0x0FFFFFFF8				0x0000000A				0x00000006			
0x20	0x00000000				0x0000F541				0x0FFFFFFF8				0x00000000			



Unallocated clusters can be identified by checking for zeroes in the FAT!

File System





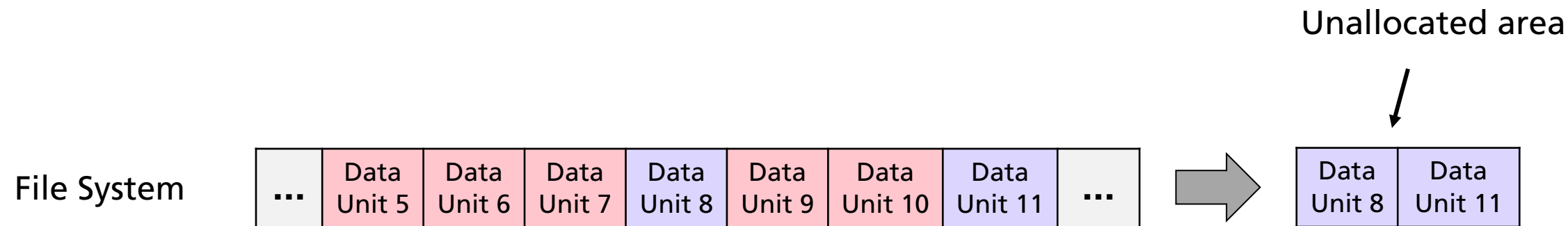
FAT file system

Analysis: Unallocated Area

- Unallocated areas may contain content of deleted files
- Can also be used to hide data since it is not actively used



byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00	0x0FFFFFFF8				0x0FFFFFFF				0x00000005				0x00000004			
0x10	0x00000007				0x0FFFFFFF8				0x0000000A				0x00000006			
0x20	0x00000000				0x0000F541				0x0FFFFFFF8				0x00000000			



The Sleuth Kit

blkls

blkls(1) - Linux man page

Name

blkls - List or output file system data units.

Synopsis

```
blkls [-aAeIsvV] [-f fstype ] [-i imgtype ] [-o imgoffset ] [-b dev sector size]  
image [images] [start-stop]
```

Description

blkls opens the named image(s) and copies file system data units (blocks). By default, **blkls** copies the contents of unallocated data blocks. **blkls** was called **dls** in TSK versions prior to 3.0.0. **blkls** was called **unrm** in TCT.

-e Copy every block, including file system metadata blocks. The output is the entire file system.

-a Display all allocated blocks (same as **-e** if **-A** is also given).

-A Display all unallocated blocks (same as **-e** if **-a** is also given). This is the default behavior.

[...]



<https://www.sleuthkit.org>

FAT file system

Name	Start	Size	Created on
file.docx	Data Unit 40	1400 Bytes	2015-07-10

Metadata



file.docx

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00	short name											attr.	res.	created time		
0x10	created date	date last accessed		first cluster (high bytes)		time of last write		date of last write		first cluster (low bytes)		file size				

Directory Entry

The Sleuth Kit

`istat`

`istat(1)` - Linux man page

Name

`istat` - Display details of a meta-data structure (i.e. inode)

Synopsis

```
istat [-B num ] [-f fstype ] [-i imgtype] [-o imgoffset] [-b dev_sector_size] [-vV] [-z zone ] [-s seconds ] image [images] inode
```

Description

`istat` displays the uid, gid, mode, size, link number, modified, accessed, changed times, and all the disk units a structure has allocated.



<https://www.sleuthkit.org>

The Sleuth Kit

istat

```
$ istat disk.dd 4
```

```
Directory Entry: 4
```

```
Allocated
```

```
File Attributes: File, Archive
```

```
Size: 1096213
```

```
Name: IMAGES~1.JPG
```

```
Directory Entry Times:
```

```
Written:      2020-11-15 18:05:06 (CET)
```

```
Accessed:    2020-11-15 00:00:00 (CET)
```

```
Created:     2020-11-15 18:05:06 (CET)
```

```
Sectors:
```

```
440 441 442 443 444 445 446 447
```

```
448 449 450 451 452 453 454 455
```

```
456 457 458 459 460 461 462 463
```

```
464 465 466 467 468 469 470 471
```

```
[...]
```



<https://www.sleuthkit.org>

The Sleuth Kit

istat

```
$ istat disk.dd 4
```

```
Directory Entry: 4
```

```
Allocated
```

```
File Attributes: File, Archive
```

```
Size: 1096213
```

```
Name: IMAGES~1.JPG
```

```
Directory Entry Times:
```

```
Written:      2020-11-15 18:05:06 (CET)
```

```
Accessed:    2020-11-15 00:00:00 (CET)
```

```
Created:     2020-11-15 18:05:06 (CET)
```

```
Sectors:
```

```
440 441 442 443 444 445 446 447
```

```
448 449 450 451 452 453 454 455
```

```
456 457 458 459 460 461 462 463
```

```
464 465 466 467 468 469 470 471
```

```
[...]
```

File Name Category

IMAGES~1.JPG



<https://www.sleuthkit.org>

The Sleuth Kit

istat

```
$ istat disk.dd 4
```

```
Directory Entry: 4
```

```
Allocated
```

```
File Attributes: File, Archive
```

```
Size: 1096213
```

```
Name: IMAGES~1.JPG
```

```
Directory Entry Times:
```

```
Written:      2020-11-15 18:05:06 (CET)
```

```
Accessed:    2020-11-15 00:00:00 (CET)
```

```
Created:     2020-11-15 18:05:06 (CET)
```

```
Sectors:
```

```
440 441 442 443 444 445 446 447
```

```
448 449 450 451 452 453 454 455
```

```
456 457 458 459 460 461 462 463
```

```
464 465 466 467 468 469 470 471
```

```
[...]
```

File Name Category

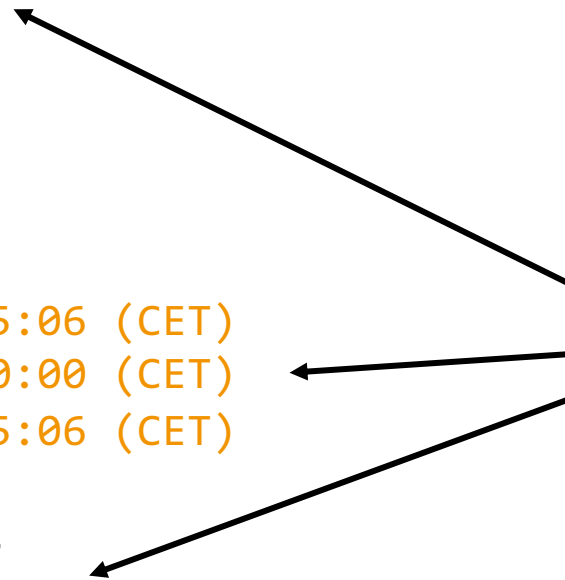
IMAGES~1.JPG

Metadata Category

Times and
Addresses



<https://www.sleuthkit.org>



The Sleuth Kit

istat

```
$ istat disk.dd 4
```

Directory Entry: 4

Allocated

File Attributes: File, Archive

Size: 1096213

Name: IMAGES~1.JPG

Directory Entry Times:

Written: 2020-11-15 18:05:06 (CET)

Accessed: 2020-11-15 00:00:00 (CET)

Created: 2020-11-15 18:05:06 (CET)

Sectors:

440 441 442 443 444 445 446 447
448 449 450 451 452 453 454 455
456 457 458 459 460 461 462 463
464 465 466 467 468 469 470 471
[...]



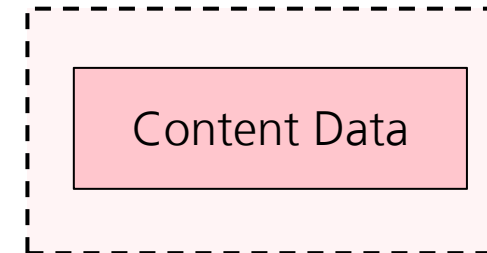
File Name Category



Metadata Category



Content Category



<https://www.sleuthkit.org>

The Sleuth Kit

`icat`

`icat(1)` - Linux man page

Name

`icat` - Output the `contents of a file` based on its inode number.

Synopsis

```
icat [-hrsvV] [-f fstype ] [-i imgtype ] [-o imgoffset ] [-b  
dev_sector_size] image [images] inode
```

Description

`icat` opens the named *image(s)* and copies the file with the specified *inode* number to standard output.



<https://www.sleuthkit.org>

The Sleuth Kit

istat

```
$ istat disk.dd 4
```

Directory Entry: 4

Allocated

File Attributes: File, Archive

Size: 1096213

Name: IMAGES~1.JPG

Directory Entry Times:

Written: 2020-11-15 18:05:06 (CET)

Accessed: 2020-11-15 00:00:00 (CET)

Created: 2020-11-15 18:05:06 (CET)

Sectors:

440 441 442 443 444 445 446 447

448 449 450 451 452 453 454 455

456 457 458 459 460 461 462 463

464 465 466 467 468 469 470 471

[...]

How are directory entries numbered?



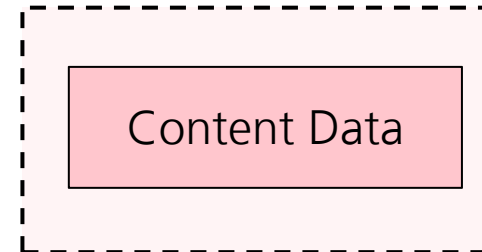
File Name Category



Metadata Category

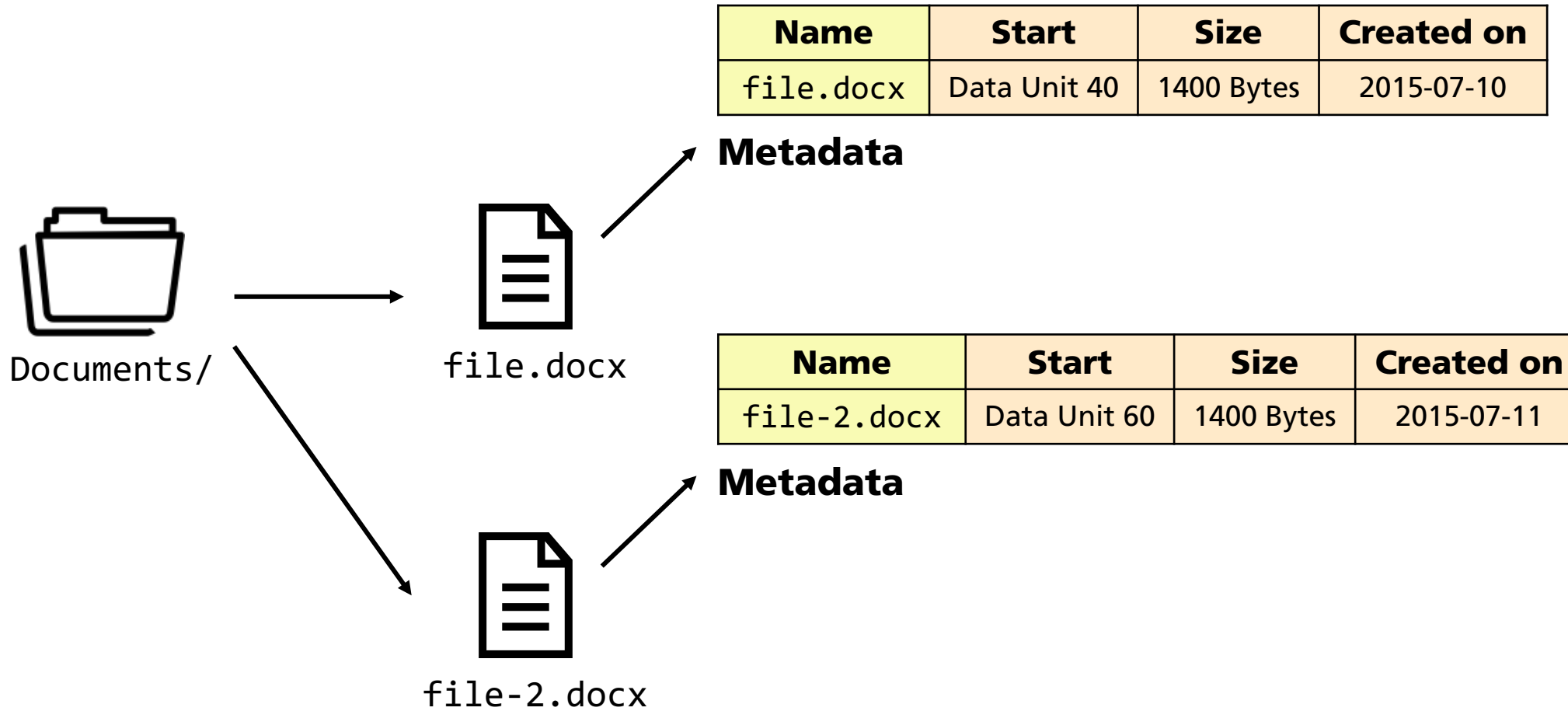


Content Category



<https://www.sleuthkit.org>

FAT file system



FAT file system

Name	Start	Size
Documents	Data Unit 20	0 Bytes

Metadata



Documents/

Directories have the corresponding attribute flag set

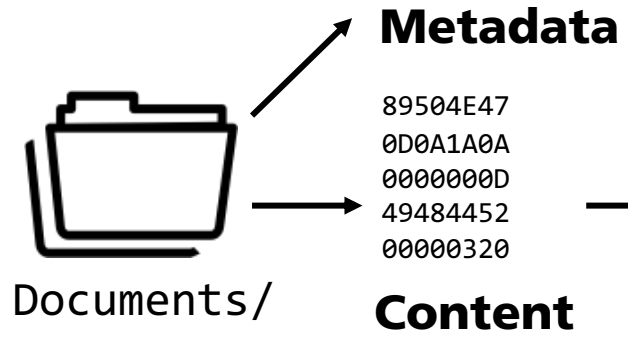
byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00	short name											attr.	res.	created time		
0x10	created date	date last accessed		first cluster (high bytes)		time of last write		date of last write		first cluster (low bytes)		0x00000000				

Directory Entry

Should be set to zero

FAT file system

Name	Start	Size
Documents	Data Unit 20	0 Bytes



Name	Start	Size
.	Data Unit 20	0
..	Data Unit 10	0
IMG10.png	Data Unit 1	2 MB
IMG11.png	Data Unit 3	500 kB
IMG12.png	Data Unit 4	500 KB
IMG13.png	Data Unit 6	1 MB
file.docx	Data Unit 40	1400 Bytes
...

The content of a directory are directory entries

FAT file system

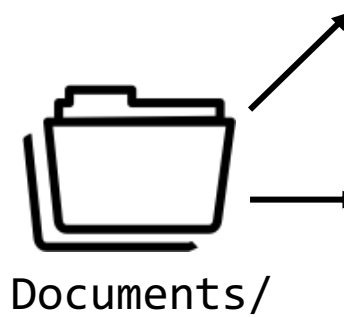
Name	Start	Size
Documents	Data Unit 20	0 Bytes

Metadata

89504E47
 0D0A1A0A
 0000000D
 49484452
 00000320

Content

Name	Start	Size
.	Data Unit 20	0
..	Data Unit 10	0
IMG10.png	Data Unit 1	2 MB
IMG11.png	Data Unit 3	500 kB
IMG12.png	Data Unit 4	500 KB
IMG13.png	Data Unit 6	1 MB
file.docx	Data Unit 40	1400 Bytes
...

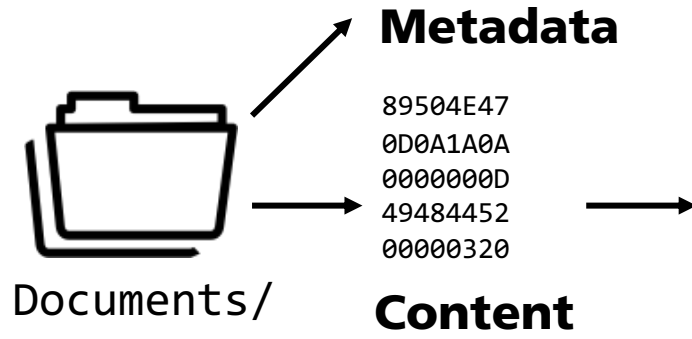


It contains directory entries for itself and its parent directory

FAT file system

Name	Start	Size
Documents	Data Unit 20	0 Bytes

- When a **new directory** is created, the assigned cluster is **wiped with zeroes**.
- The **creation timestamp** of the **.** and **..** entries should match the **creation timestamp** within the parent directory



Name	Start	Size
.	Data Unit 20	0
..	Data Unit 10	0
IMG10.png	Data Unit 1	2 MB
IMG11.png	Data Unit 3	500 kB
IMG12.png	Data Unit 4	500 KB
IMG13.png	Data Unit 6	1 MB
file.docx	Data Unit 40	1400 Bytes
Secret	Data Unit 90	0
...

Cluster 20

Cluster 90

Name	Start	Size
.	Cluster 90	0
..	Cluster 20	0
top.pdf	Cluster 190	1 MB
...

FAT file system

Metadata Addresses



Documents/

File System



FAT file system

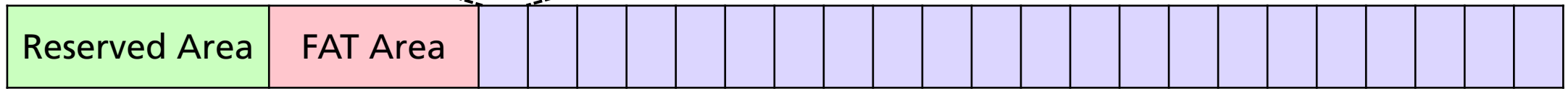
Metadata Addresses



Documents/

Name	Start	Size	Metadata Address
IMG10.png	Cluster 40	1400 Bytes	3
IMG11.png			4
IMG12.png			5
IMG13.png			6
IMG14.png			7
...

File System



Sector 520

(first sector within data area)

FAT file system

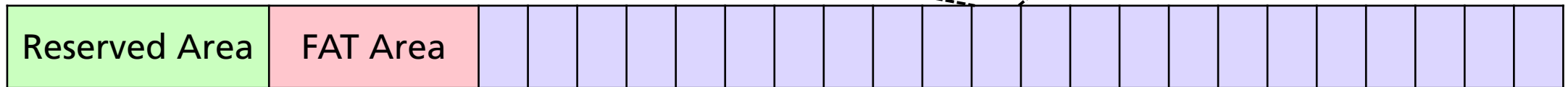
Metadata Addresses



Documents/

Name	Start	Size	Metadata Address
			163
			164
			165
			166
			167
...

File System



Sector 530

FAT file system

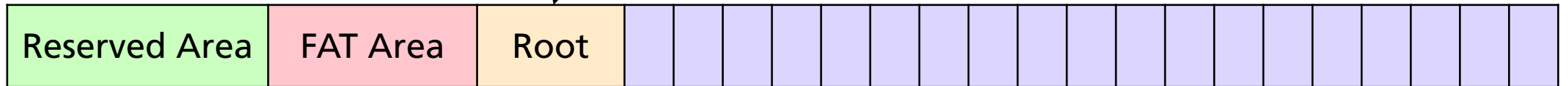
Metadata Addresses



Documents/

Region used for the root directory entries (FAT12 and FAT16)

File System



The Sleuth Kit

`fls`

`fls(1)` - Linux man page

Name

`fls` - List `file and directory names` in a disk image.

Synopsis

```
fls [-adDFlpruvV] [-m mnt ] [-z zone ] [-f fstype ] [-s seconds ] [-i imgtype ]  
[-o imgoffset ] [-b dev_sector_size] image [images] [ inode ]
```

Description

`fls` lists the files and directory names in the *image* and can display file names of recently deleted files for the directory using the given *inode*. If the *inode* argument is not given, the *inode* value for the root directory is used. For example, on an NTFS file system it would be 5 and on a Ext3 file system it would be 2.



<https://www.sleuthkit.org>

The Sleuth Kit

`fls`

```
$ fls -r disk.dd
```

```
r/r 4:  suspicios.jpg
```

```
d/d 6:  videos
```

```
+ r/r 43014:  vid1.mp4
```

```
+ r/r 43016:  vid2.mp4
```

```
r/r * 7:  _UN-56~1.JPG
```

```
v/v 3270339:  $MBR
```

```
v/v 3270340:  $FAT1
```

```
v/v 3270341:  $FAT2
```

```
V/V 3270342:  $OrphanFiles
```



<https://www.sleuthkit.org>

The Sleuth Kit

fls

```
$ fls -r disk.dd
```

```
r/r 4: suspicios.jpg
```

```
d/d 6: videos
```

```
+ r/r 43014: vid1.mp4
```

```
+ r/r 43016: vid2.mp4
```

```
r/r * 7: _UN-56~1.JPG
```

```
v/v 3270339: $MBR
```

```
v/v 3270340: $FAT1
```

```
v/v 3270341: $FAT2
```

```
V/V 3270342: $OrphanFiles
```

Metadata (inode) addresses



<https://www.sleuthkit.org>

The Sleuth Kit

fls

```
$ fls -r disk.dd
```

```
r/r 4: suspicios.jpg
```

```
d/d 6: videos
```

```
+ r/r 43014: vid1.mp4
```

```
+ r/r 43016: vid2.mp4
```

```
r/r * 7: _UN-56~1.JPG
```

```
v/v 3270339: $MBR
```

```
v/v 3270340: $FAT1
```

```
v/v 3270341: $FAT2
```

```
V/V 3270342: $OrphanFiles
```

Deleted files are indicated
by an asterisk



<https://www.sleuthkit.org>

FAT file system

Scenario: File Deletion

Name	Start	Size
file.docx	Cluster 40	1400 Bytes
NYC.png	Cluster 60	5000 Bytes
IMG20.png	Cluster 100	7000 Bytes
videos	Cluster 10	0
documents	Cluster 11	0
...

Cluster 140

	...
58	0
59	0
60	62
61	EOF
62	63
63	EOF
	...

FAT

FAT file system

Scenario: File Deletion

Name	Start	Size
file.docx	Cluster 40	1400 Bytes
0xe5YC.png	Cluster 60	5000 Bytes
IMG20.png	Cluster 100	7000 Bytes
videos	Cluster 10	0
documents	Cluster 11	0
...

Cluster 140

	...
58	0
59	0
60	62
61	EOF
62	63
63	EOF
	...

FAT

- After deletion, the metadata entry is marked as unallocated by setting the first byte to 0xe5

FAT file system

Scenario: File Deletion

Name	Start	Size
file.docx	Cluster 40	1400 Bytes
0xe5YC.png	Cluster 60	5000 Bytes
IMG20.png	Cluster 100	7000 Bytes
videos	Cluster 10	0
documents	Cluster 11	0
...

Cluster 140

	...
58	0
59	0
60	0
61	EOF
62	0
63	0
	...

FAT

- After deletion, the metadata entry is marked as unallocated by setting the first byte to `0xe5`
- The corresponding clusters are marked as unallocated by setting their entry in the FAT to `0`

„File deletion is *OS-specific* and may differ between *implementations*. Most implementations of FAT file systems will delete a file by *setting the directory entry to an unallocated state* (which overwrites the first letter of the name) and *sets the table entries for the file's clusters to 0*. In general, the starting cluster and the size of the file are not wiped from the directory entry. Again, this is *OS-specific* and an OS may choose to wipe those fields when a file is deleted.“

<http://www.sleuthkit.org/informer/sleuthkit-informer-14.txt>

FAT file system

Analysis: Metadata-based File Recovery

Name	Start	Size
file.docx	Cluster 40	1400 Bytes
0xe5YC.png	Cluster 60	5000 Bytes
IMG20.png	Cluster 100	7000 Bytes
videos	Cluster 10	0
documents	Cluster 11	0
...

Cluster 140

	...
58	0
59	0
60	0
61	EOF
62	0
63	0
	...

FAT

- After deletion, the metadata entry is marked as unallocated by setting the first byte to `0xe5`
- The corresponding clusters are marked as unallocated by setting their entry in the FAT to `0`

- Unallocated metadata entries can still be found, e.g. by searching for a certain signature (such as `0xe5`) or by simply iterating over a data unit
- Referenced data units may have already been overwritten without you noticing
- Even worse: If metadata entries don't store all content addresses, it may be impossible to recover a file



FAT file system

Analysis: Metadata-based File Recovery

Name	Start	Size
file.docx	Cluster 40	1400 Bytes
0xe5YC.png	Cluster 60	5000 Bytes
IMG20.png	Cluster 100	7000 Bytes
videos	Cluster 10	0
documents	Cluster 11	0
LA.png	Cluster 60	6000 Bytes
...

Cluster 140

	...
58	0
59	0
60	140
61	EOF
62	0
63	0
	...

FAT

- After deletion, the metadata entry is marked as unallocated by setting the first byte to `0xe5`
- The corresponding clusters are marked as unallocated by setting their entry in the FAT to `0`

- Unallocated metadata entries can still be found, e.g. by searching for a certain signature (such as `0xe5`) or by simply iterating over a data unit
- Referenced data units may have already been overwritten without you noticing
- Even worse: If metadata entries don't store all content addresses, it may be impossible to recover a file



The Sleuth Kit

fls & icat

```
$ fls -r disk.dd
```

```
r/r 4: suspicios.jpg
```

```
d/d 6: videos
```

```
+ r/r 43014: vid1.mp4
```

```
+ r/r 43016: vid2.mp4
```

```
r/r * 7: _UN-56~1.JPG
```

```
v/v 3270339: $MBR
```

```
v/v 3270340: $FAT1
```

```
v/v 3270341: $FAT2
```

```
V/V 3270342: $OrphanFiles
```

Deleted files are indicated
by an asterisk



<https://www.sleuthkit.org>

```
$ icat disk.dd 7 > recovered.JPG
```

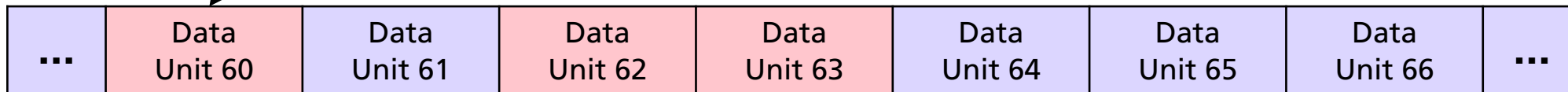
The Sleuth Kit

FAT File Recovery



<https://www.sleuthkit.org>

Name	Start	Size
NYC.png	Cluster 60	5000 Bytes



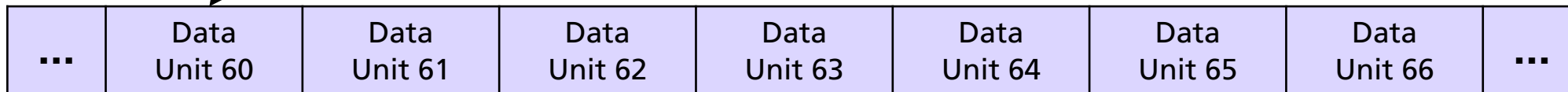
The Sleuth Kit

FAT File Recovery



<https://www.sleuthkit.org>

Name	Start	Size
0xe5YC.png	Cluster 60	5000 Bytes



The Sleuth Kit

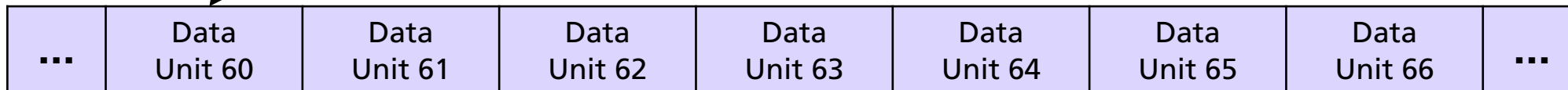
FAT File Recovery

- For FAT file recovery TSK uses the **start** and **size fields** of **unallocated** directory entries
- If the starting cluster is **allocated**, it will **not recover any data**



<https://www.sleuthkit.org>

Name	Start	Size
0xe5YC.png	Cluster 60	5000 Bytes



The Sleuth Kit

FAT File Recovery

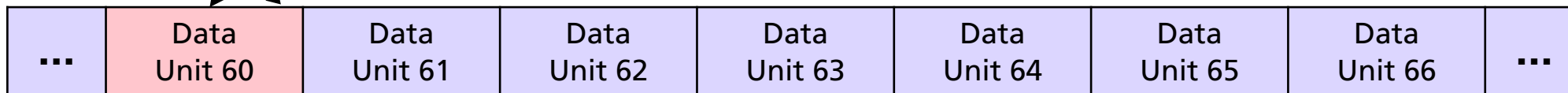
- For FAT file recovery TSK uses the **start** and **size fields** of **unallocated** directory entries
- If the starting cluster is **allocated**, it will **not recover any data**



<https://www.sleuthkit.org>

Name	Start	Size
0xe5YC.png	Cluster 60	5000 Bytes

Name	Start	Size
LA.png	Cluster 60	6000 Bytes



The Sleuth Kit

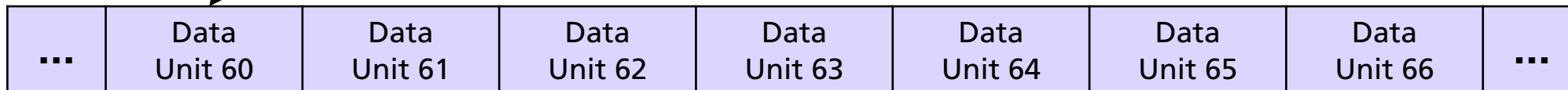
FAT File Recovery

- For FAT file recovery TSK uses the **start** and **size fields** of **unallocated** directory entries
- If the starting cluster is **allocated**, it will **not recover any data**
- If the starting cluster is **unallocated**, it will extract **consecutive** and **unallocated** clusters until the **final file size** is reached



<https://www.sleuthkit.org>

Name	Start	Size
0xe5YC.png	Cluster 60	5000 Bytes



The Sleuth Kit

FAT File Recovery

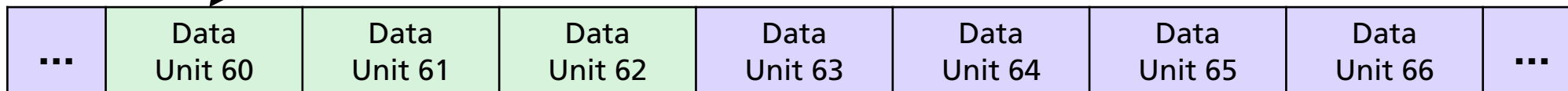
- For FAT file recovery TSK uses the **start** and **size fields** of **unallocated** directory entries
- If the starting cluster is **allocated**, it will **not recover any data**
- If the starting cluster is **unallocated**, it will extract **consecutive** and **unallocated** clusters until the **final file size** is reached



<https://www.sleuthkit.org>

Name	Start	Size
0xe5YC.png	Cluster 60	5000 Bytes

Wrong extraction for file
NCY.png



The Sleuth Kit

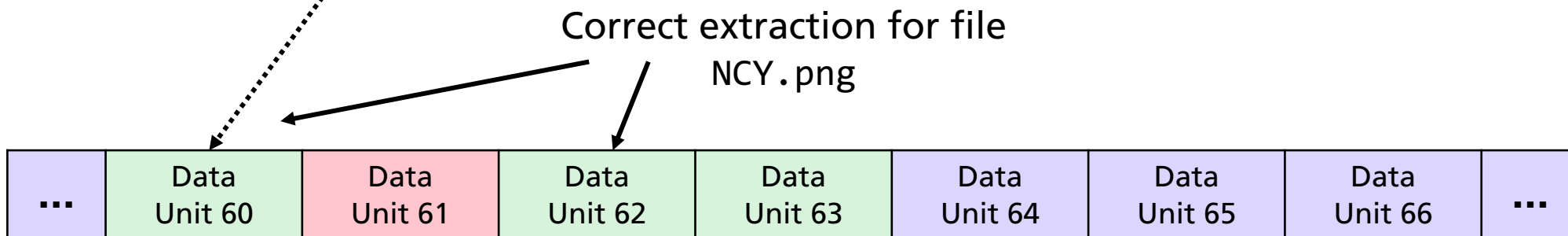
FAT File Recovery

- For FAT file recovery TSK uses the **start** and **size fields** of **unallocated** directory entries
- If the starting cluster is **allocated**, it will **not recover any data**
- If the starting cluster is **unallocated**, it will extract **consecutive** and **unallocated** clusters until the **final file size** is reached



<https://www.sleuthkit.org>

Name	Start	Size
0xe5YC.png	Cluster 60	5000 Bytes



FAT file system

Scenario: Directory Deletion

Name	Start	Size
file.docx	Cluster 40	1400 Bytes
0xe5YC.png	Cluster 60	5000 Bytes
IMG20.png	Cluster 100	7000 Bytes
videos	Cluster 10	0
documents	Cluster 11	0
LA.png	Cluster 60	6000 Bytes
...

Cluster 140



Name	Start	Size
.	Cluster 10	0
..	Cluster 110	0
movie.mp4	Cluster 200	10 MB
...

Cluster 10

FAT file system

Scenario: Directory Deletion

Name	Start	Size
file.docx	Cluster 40	1400 Bytes
0xe5YC.png	Cluster 60	5000 Bytes
IMG20.png	Cluster 100	7000 Bytes
0xe5ideos	Cluster 10	0
documents	Cluster 11	0
LA.png	Cluster 60	6000 Bytes
...

Cluster 140

Name	Start	Size
.	Cluster 10	0
..	Cluster 110	0
0xe5ovie.mp4	Cluster 200	10 MB
...

Cluster 10

- Directory entries of **deleted directories** are also **unallocated** except for the **first two!**
- **Clusters** belonging to **directories** can easily be identified by the **.** and **..** directory entries at the **beginning**



FAT file system

Scenario: Directory Deletion

Name	Start	Size
file.docx	Cluster 40	1400 Bytes
0xe5YC.png	Cluster 60	5000 Bytes
IMG20.png	Cluster 100	7000 Bytes
music	Cluster 200	0
documents	Cluster 11	0
LA.png	Cluster 60	6000 Bytes
...

Cluster 140

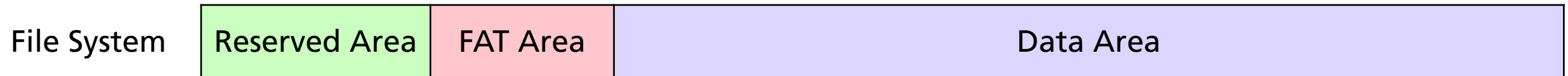
Name	Start	Size
.	Cluster 10	0
..	Cluster 110	0
0xe5ovie.mp4	Cluster 200	10 MB
...

Cluster 10

Files like these are
called "Orphan Files"

https://wiki.sleuthkit.org/index.php?title=Orphan_Files

FAT file system



FAT file system

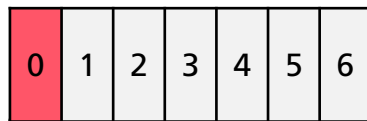
Reserved Area



FAT file system

Reserved Area

First sector is the **boot sector**



File System



FAT file system

Reserved Area

For FAT32:
File System Information Sector (FSINFO)



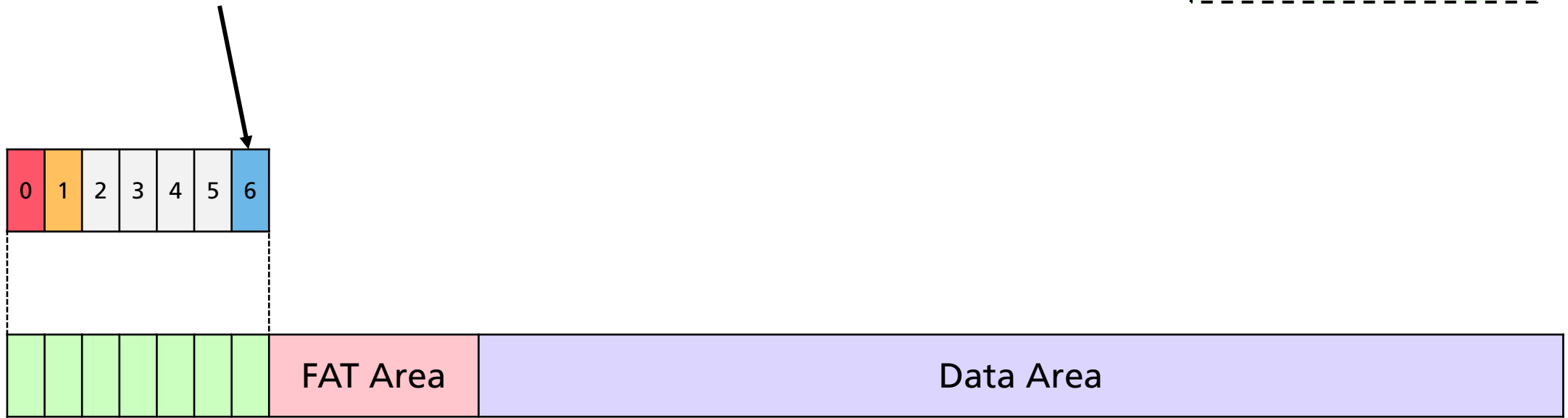
FAT file system

Reserved Area

File System Category

Layout and Size Information

FAT32 also contains a backup boot sector



FAT file system

Boot Sector

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f	
0x00	jump to boot code			OEM name in ASCII													
0x10	BIOS Parameter Block																
0x20						drive no.	res.	boot sig.	volume serial number				volume label				
0x30	in ASCII						file system type label in ASCII										
...	not used / boot code																
0x1f																0xAA55	

Signature value at offset 510, which is usually also the end of the sector

FAT file system

Boot Sector

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f		
0x00	jump to boot code			OEM name in ASCII														
0x10	BIOS Parameter Block																	
0x20						drive no.	res.	boot sig.	volume serial number				volume label					
0x30	in ASCII							file system type label in ASCII										
...	not used / boot code																	
0x1f																0xAA55		



Signature values are a great reference to look for in case of missing or damaged data (e.g. missing partition tables).

0xAA55

Signature Detection

sigfind

sigfind(1) - Linux man page

Name

sigfind - Find a binary signature in a file

Synopsis

sigfind [-b bsize] [-o offset] [-t template] [-lV] [hex signature] file

Description

sigfind searches through a file and looks for the `hex_signature` at a given offset. This can be used to search for lost `boot sectors`, superblocks, and partition tables.



<https://www.sleuthkit.org>

Usually, the start of a file system can be inferred from the partition layout



Signature Detection

sigfind

sigfind(1) - Linux man page

Name

sigfind - Find a binary signature in a file

Synopsis

```
sigfind [-b bsize ] [-o offset ] [-t template ] [-lV] [ hex signature ] file
```

Description

sigfind searches through a file and looks for the `hex_signature` at a given offset. This can be used to search for lost `boot sectors`, superblocks, and partition tables.

```
$ sigfind -o 510 -l AA55 disk.dd
```



<https://www.sleuthkit.org>

Signature Detection

sigfind

sigfind(1) - Linux man page

Name

sigfind - Find a binary signature in a file

Synopsis

sigfind [-b bsize] [-o offset] [-t template] [-lV] [hex signature] file

Description

sigfind searches through a file and looks for the hex_signature at a given offset. This can be used to search for lost boot sectors, superblocks, and partition tables.



<https://www.sleuthkit.org>

```
$ sigfind -o 510 -l AA55 disk.dd
```



Signature Detection

gpart

gpart(8) - Linux man page

Name

gpart - guess PC-type hard disk partitions

Synopsis

gpart [options] device

Options: [-b <backup MBR>][-C c,h,s][-c][-d][-E][-e][-f] [-g][-h][-i]
[-K <last-sector>][-k <# of sectors>] [-L] [-l <log file>]
[-n <increment>] [-q][-s <sector-size>] [-t <module-name>][-V][-v]
[-W <device>][-w <module-name, weight>]

Description

gpart tries to guess which partitions are on a hard disk. If the primary partition table has been lost, overwritten or destroyed the partitions still exist on the disk but the operating system cannot access them.

FAT file system

Boot Sector

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00	jump to boot code			OEM name in ASCII												
0x10	BIOS Parameter Block															
0x20					drive no.	res.	boot sig.	volume serial number				volume label				
0x30	in ASCII						file system type label in ASCII									
...	not used / boot code															
0x1f																0xAA55

FAT file system

Boot Sector

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00												BIOS Parameter Block				
0x10																
0x20																
0x30																
...																
0x1f																

FAT file system

BIOS Parameter Block

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f	
0x00												sector size					
0x10	Size of a sector in bytes (usually 512, 1024, 2048 or 4096)																
0x20																	
0x30																	
...																	
0x1f																	

mkfs.fat - create an MS-DOS FAT filesystem

-S *LOGICAL-SECTOR-SIZE* Specify the number of bytes per logical sector. Must be a power of 2 and greater than or equal to 512, i.e. 512, 1024, 2048, 4096, 8192, 16384, or 32768. Values larger than 4096 are not conforming to the FAT file system specification and may not work everywhere.

<https://www.man7.org/linux/man-pages/man8/mkfs.fat.8.html>

FAT file system

BIOS Parameter Block

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f	
0x00												sector size	clust. size				
0x10																	
0x20																	
0x30																	
...																	
0x1f																	

Sectors per cluster.
Value must be a power of 2 (i.e. 1, 2, 4, 8, 16, 32, 64, 128)

FAT file system

BIOS Parameter Block

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00												sector size	clust. size	reserved area size		
0x10																
0x20																
0x30																
...																
0x1f																

Number of sectors in the reserved area

FAT file system

BIOS Parameter Block

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00												sector size	clust. size	reserved area size		
0x10	# of FATs	max. # of root entries	total sectors	media type	FAT size in sectors	geometry information						hidden sectors				
0x20	total sectors (4 bytes)															
0x30																
...																
0x1f																

FAT file system

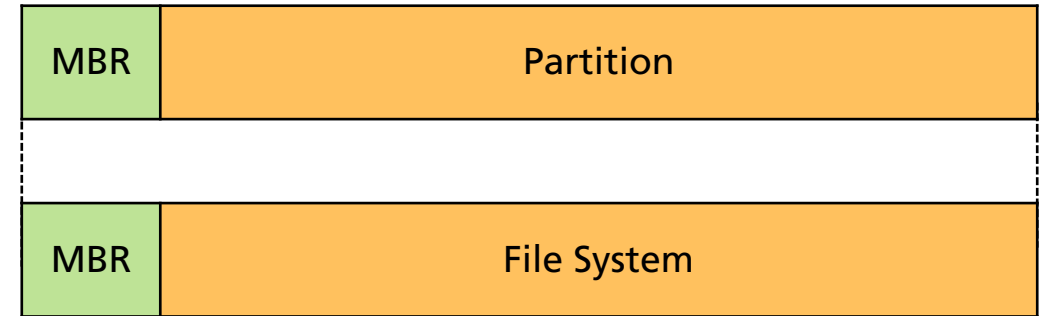
BIOS Parameter Block

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00												sector size	clust. size	reserved area size		
0x10	# of FATs	max. # of root entries	total sectors	media type	FAT size in sectors	geometry information						hidden sectors				
0x20	total sectors (4 bytes)															
0x30																
...																
0x1f																



FAT file system

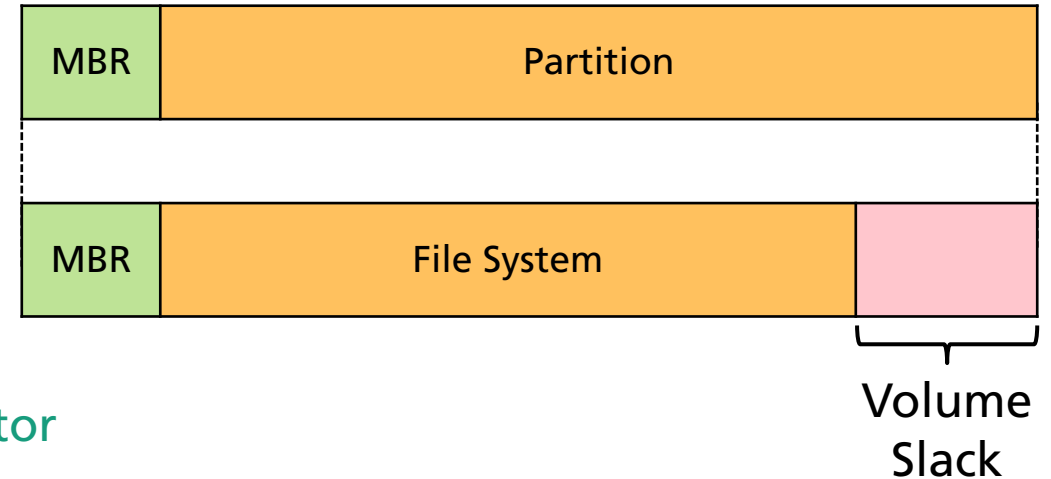
Analysis: Volume Slack



FAT file system

Analysis: Volume Slack

- **Volume Slack** describes the space between the end of a file system and the end of the volume it is stored on
- It can be detected by comparing the total size of a file system with the size of the volume (e.g. partition)
 - For **FAT**, this information is found in the **boot sector**
- Creating volume slack is easy, since only one parameter has to be modified



FAT file system

BIOS Parameter Block

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00												sector size	clust. size	reserved area size		
0x10	# of FATs	max. # of root entries	total sectors	media type	FAT size in sectors	geometry information						hidden sectors				
0x20	total sectors (4 bytes)															
0x30																
...																
0x1f																

FAT file system

BIOS Parameter Block

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00												sector size	clust. size	reserved area size		
0x10	# of FATs	max. # of root entries	total sectors	media type	FAT size in sectors	geometry information						hidden sectors				
0x20	total sectors (4 bytes)															
0x30																
...																
0x1f																

Number of sectors before the start of the partition the file system is using (beginning of the partition)




Number of sectors before the start of the partition the file system is using (beginning of the partition)



FAT file system

BIOS Parameter Block

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00												sector size	clust. size	reserved area size		
0x10	# of FATs	max. # of root entries	total sectors	media type	FAT size in sectors	geometry information						hidden sectors				
0x20	total sectors (4 bytes)															
0x30	<div style="text-align: right;"> File System Category  </div>															
...																
0x1f	<div style="border: 1px dashed black; padding: 10px; text-align: center;"> <div style="border: 1px solid black; padding: 5px; display: inline-block;">Layout and Size Information</div> </div>															

File System





FAT file system

BIOS Parameter Block (FAT32)

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00																
0x10	same as for FAT12/FAT16															
0x20																
0x30																
0x40																
0x50																
...																

FAT file system

BIOS Parameter Block (FAT32)

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f	
0x00																	
0x10	same as for FAT12/FAT16																
0x20						FAT size in sectors				FAT usage flags		version		root directory cluster			
0x30	FSINFO sector		Backup boot sector		reserved												
0x40																	
0x50																	
...																	

FAT file system

FSINFO

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00	0x41615252															
...																
0x1e																
0x1f																

FSINFO sector signature



FAT file system

FSINFO

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00	0x41615252															
...	not used															
0x1e																
0x1f																

FAT file system

FSINFO

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00	0x41615252															
...	not used															
0x1e					0x61417272											
0x1f																

Another signature

FAT file system

FSINFO

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00	0x41615252															
...	not used															
0x1e					0x61417272				number of free clusters							
0x1f																

Last number of free clusters (which must not be correct).
 0xFFFFFFFF indicates an unknown value.



FAT file system

FSINFO

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00	0x41615252															
...	not used															
0x1e					0x61417272				number of free clusters				next free cluster			
0x1f																

Usually set to the last allocated cluster, indicating from where to search for other free clusters.





FAT file system

FSINFO

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00	0x41615252															
...	not used															
0x1e					0x61417272				number of free clusters				next free cluster			
0x1f	reserved															



FAT file system

FSINFO

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00	0x41615252															
...	not used															
0x1e					0x61417272				number of free clusters				next free cluster			
0x1f	reserved												0xAA550000			

↑
Another signature

The Sleuth Kit

`fsstat`

`fsstat(1)` - Linux man page

Name

`fsstat` - Display `general details of a file system`

Synopsis

```
fsstat [-f fstype ] [-i imgtype] [-o imgoffset] [-b dev_sector_size] [-tvV] image  
[images]
```

Description

`fsstat` displays the details associated with a file system. The output of this command is `file system specific`. At a minimum, the range of meta-data values (inode numbers) and content units (blocks or clusters) are given. Also given are details from the Super Block, such as mount times and and features. For file systems that use groups (FFS and EXT2FS), the layout of each group is listed. For a FAT file system, the FAT table is displayed in a condensed format. Note that the data is in sectors and not in clusters.



<https://www.sleuthkit.org>



The Sleuth Kit

fsstat

```
$ fsstat disk.dd -o 2048
```

FILE SYSTEM INFORMATION

File System Type: FAT16

OEM Name: mkfs.fat

Volume ID: 0x3075b80d

Volume Label (Boot Sector): NO NAME

Volume Label (Root Directory):

File System Type Label: FAT16

Sectors before file system: 0

File System Layout (in sectors)

Total Range: 0 - 202751

* Reserved: 0 - 3

** Boot Sector: 0

* FAT 0: 4 - 203

* FAT 1: 204 - 403

* Data Area: 404 - 202751

** Root Directory: 404 - 435

** Cluster Area: 436 - 202751

FAT file system

Boot Sector

byte offset	0x00	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x0a	0x0b	0x0c	0x0d	0x0e	0x0f
0x00	jump to boot code		OEM name in ASCII													
0x10	BIOS Parameter Block															
0x20	drive no.		res.		boot sig.		volume serial number				volume label					
0x30	in ASCII						file system type label in ASCII									
...	not used / boot code															
0x1f																0xAA55

METADATA INFORMATION

Range: 2 - 3237574

Root Directory: 2

CONTENT INFORMATION

Sector Size: 512

Cluster Size: 2048

Total Cluster Range: 2 - 50580

FAT CONTENTS (in sectors)



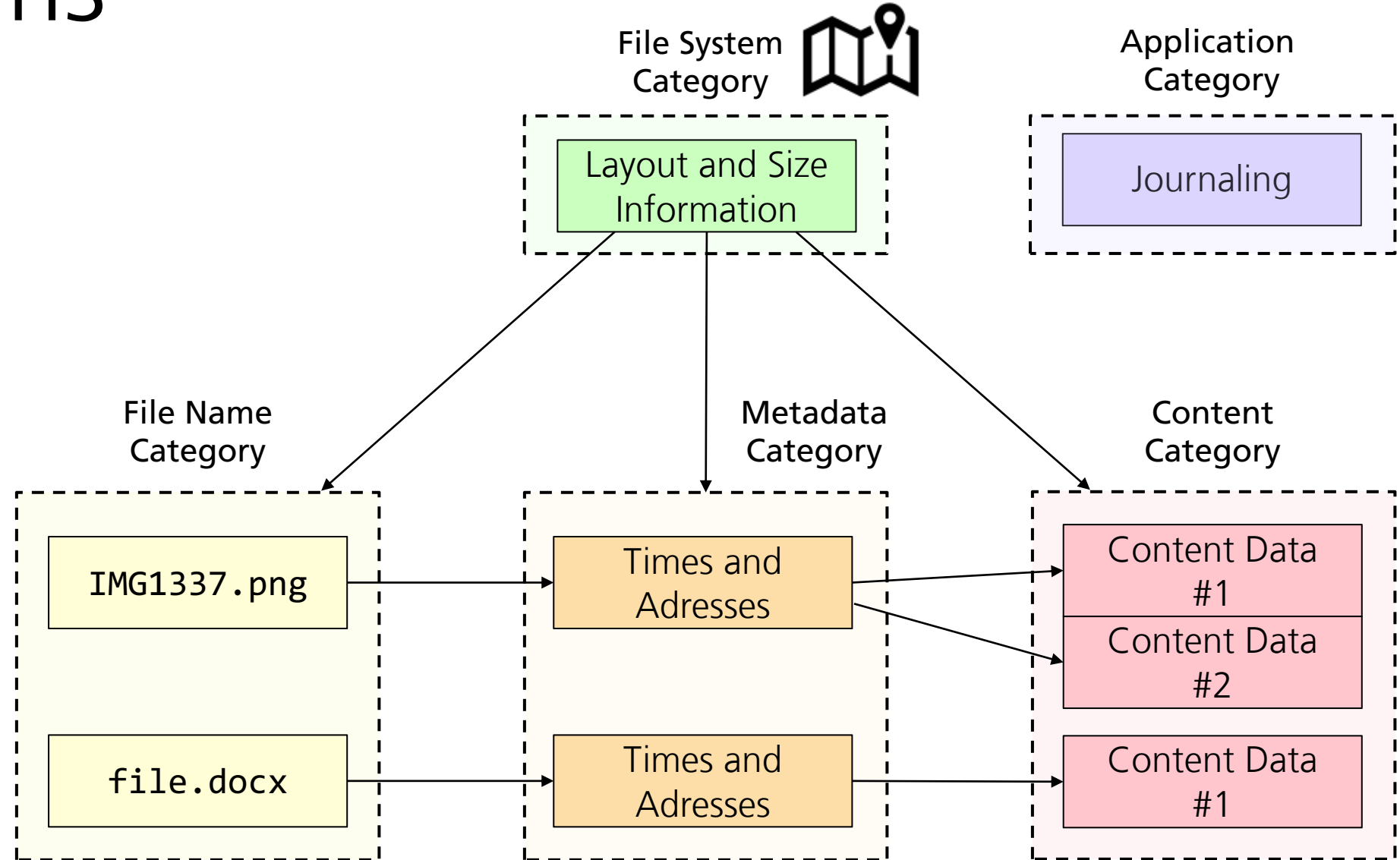
<https://www.sleuthkit.org>

File System Category



Layout and Size Information

File Systems



Dr. Brian Carrier

File Systems

- File system analysis including its **results** is **completely dependent** on the file system
- In most cases **new file systems** require **new tools** in order to achieve the most of it
- We have covered an important example, but there are more...

- **ADFS** – Acorn's Advanced Disc filing system, successor to **DFS**.
- **AdvFS** – Advanced File System, designed by **Digital Equipment Corporation** for their Digital UNIX (now **Tru64 UNIX**) operating system.
- **APFS** – Apple File System is a next-generation file system for Apple products.
- **AthFS** – AtheOS File System, a **64-bit journaled** filesystem now used by **Syllable**. Also called AFS.
- **BFS** – the Boot File System used on System V release 4.0 and UnixWare.
- **BFS** – the Be File System used on **BeOS**, occasionally misnamed as BeFS. Open source implementation called OpenBFS is used by the **Haiku** operating system.
- **Btrfs** – is a **copy-on-write** file system for **Linux** announced by **Oracle** in 2007 and published under the **GNU General Public License** (GPL).
- **CFS** – The Cluster File System from Veritas, a Symantec company. It is the parallel access version of VxFS.
- **CP/M** file system – Native filesystem used in the CP/M (Control Program for Microcomputers) operating system which was first released in 1974.
- **DOS 3.x** – Original floppy operating system and file system developed for the **Apple II**.
- **Extent File System** (EFS) – an older block filing system under **IRIX**.
- **ext** – Extended file system, designed for **Linux** systems.
- **ext2** – Second extended file system, designed for **Linux** systems.
- **ext3** – A journaled form of ext2.
- **ext4** – A follow up for **ext3** and also a **journaled** filesystem with support for **extents**.
- **ext3cow** – A **versioning file system** form of ext3.
- **FAT** – File Allocation Table, initially used on **DOS** and **Microsoft Windows** and now widely used for portable USB storage and some other devices; **FAT12**, **FAT16** and **FAT32** for 12-, 16- and 32-bit table depths.
 - **VFAT** – Optional layer on **Microsoft Windows** FAT system to allow long (up to 255 character) filenames instead of only the **8.3 filenames** allowed in the plain FAT filesystem.
 - **FATX** – A modified version of **Microsoft Windows** FAT system that is used on the original **Xbox** console.

- **FFS (Amiga)** – Fast File System, used on **Amiga** systems. This FS has evolved over time. Now counts FFS1, FFS Intl, FFS DCache, FFS2.
- **FFS** – Fast File System, used on ***BSD** systems
- **Fossil** – **Plan 9 from Bell Labs** snapshot archival file system.
- **Files-11** – **OpenVMS** file system; also used on some **PDP-11** systems; supports record-oriented files
- **Flex machine** file system
- **HAMMER** – clustered **DragonFly BSD** filesystem, production-ready since DragonFly 2.2 (2009)^{[1][2]}
- **HAMMER2** – recommended as the default root filesystem in DragonFly since 5.2 release in 2018^{[3][4][5]}
- **HFS** – Hierarchical File System in IBM's **z/OS**; not to be confused with Apple's HFS. HFS is still supported but IBM's stated direction is **zFS**.
- **HFS** – Hierarchical File System, in use until HFS+ was introduced on Mac OS 8.1. Also known as Mac OS Standard format. Successor to Macintosh File System (MFS) & predecessor to HFS+; not to be confused with IBM's HFS provided with **z/OS**
- **HFS+** – Updated version of Apple's HFS, Hierarchical File System, supported on Mac OS 8.1 & above, including macOS. Supports file system journaling, enabling recovery of data after a system crash. Also referred to as 'Mac OS Extended format or HFS Plus
- **HPFS** – High Performance File System, used on **OS/2**
- **HTFS** – High Throughput Filesystem, used on **SCO OpenServer**
- **ISO 9660** – Used on **CD-ROM** and **DVD-ROM** discs (**Rock Ridge** and **Joliet** are extensions to this)
- **JFS** – **IBM Journaling** file system, provided in **Linux**, **OS/2**, and **AIX**. Supports **extents**.
- **LFS** – 4.4BSD implementation of a **log-structured file system**
- **MFS** – Macintosh File System, used on early **Classic Mac OS** systems. Succeeded by Hierarchical File System (HFS).
- **Next3** – A form of **ext3** with snapshots support.^[6]
- **MFS** – Tivo's Media File System, a proprietary fault tolerant format used on **TiVo** hard drives for real time recording from live TV.
- **Minix file system** – Used on **Minix** systems

https://en.wikipedia.org/wiki/List_of_file_systems (part of)



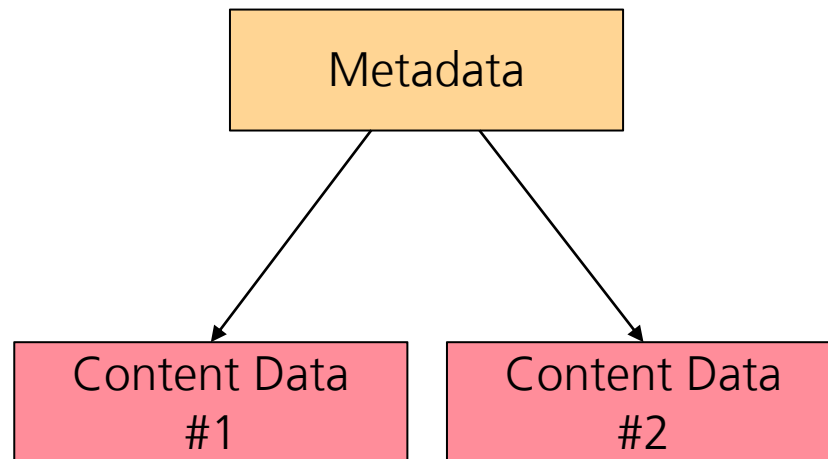
Fancy File System Techniques



Copy-on-write

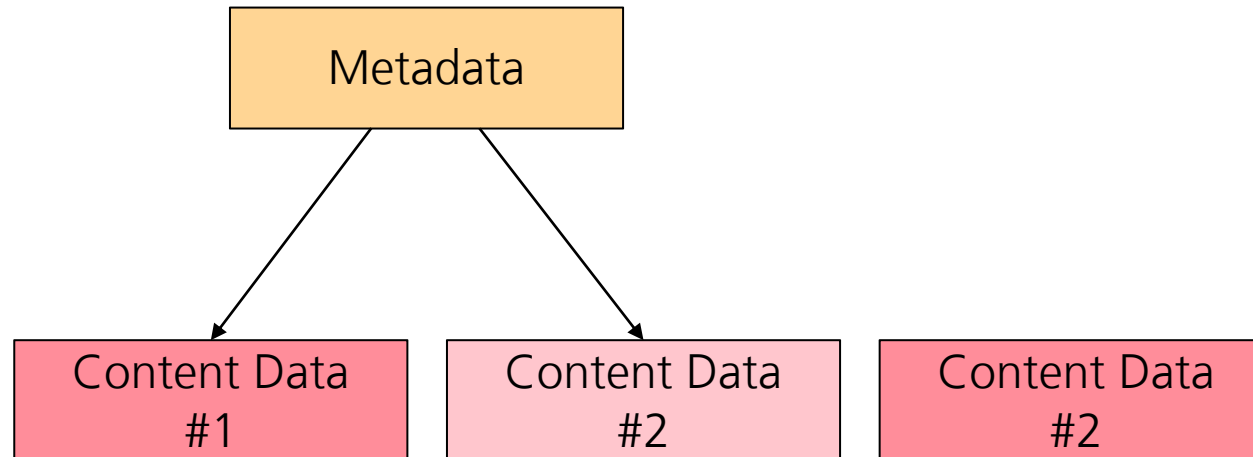
Copy-on-write

- Technique to ensure a file system stays in a **consistent state**



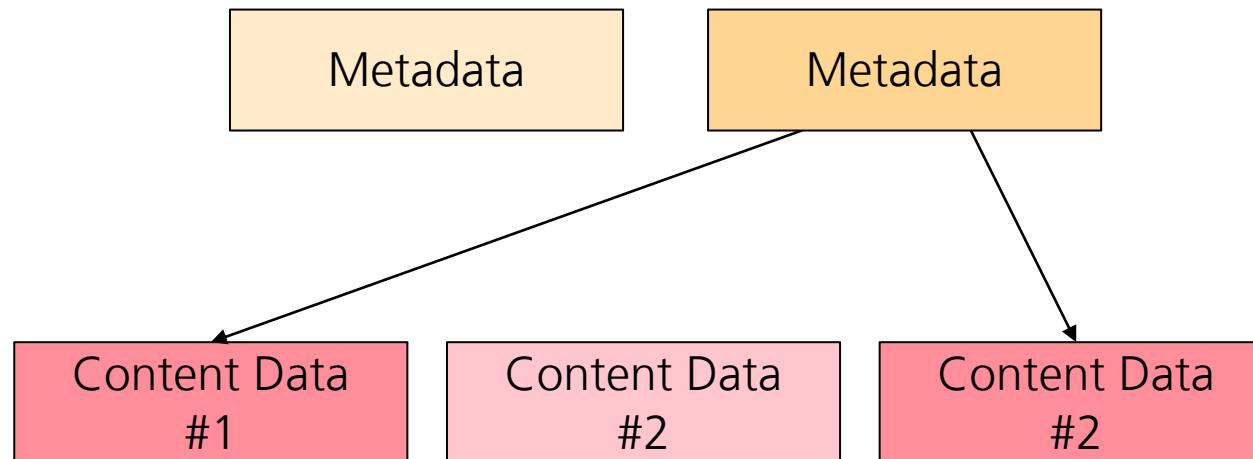
Copy-on-write

- Technique to ensure a file system stays in a **consistent state**
 - Content in the second data block is **modified**, so it is completely rewritten (**copied**) to a different area



Copy-on-write

- Technique to ensure a file system stays in a **consistent state**
 - Same applies to **metadata structures** which are updated during this process (e.g. timestamps, addresses)

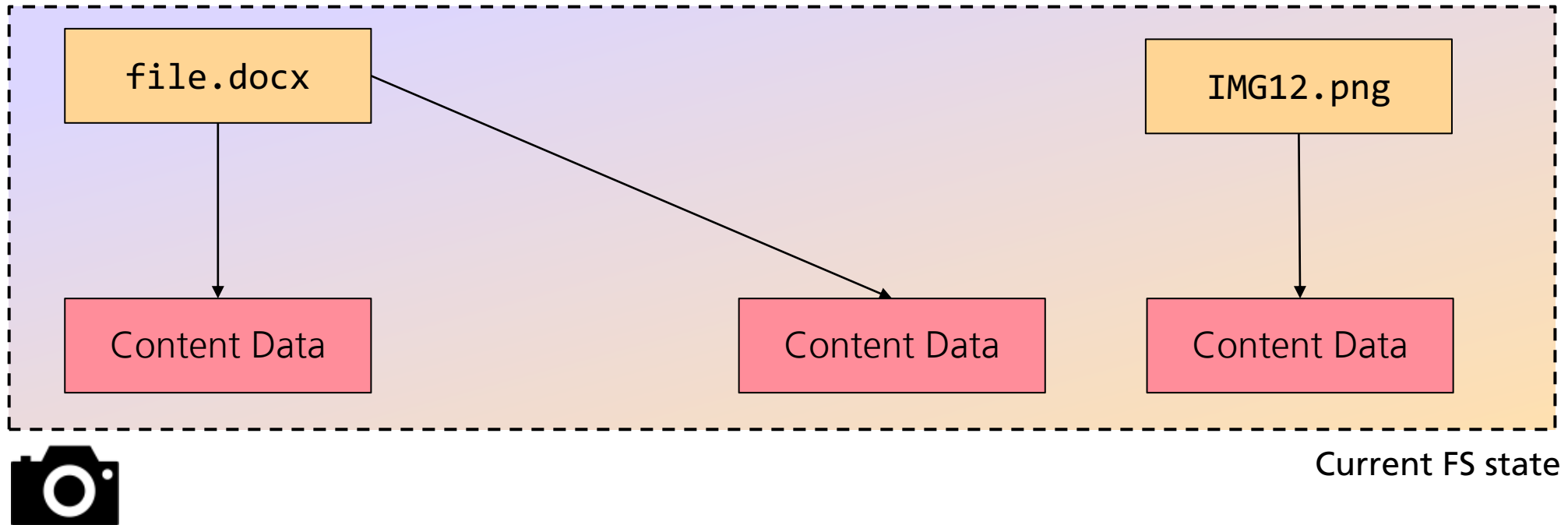


Results in a lot of artifacts!



Copy-on-write

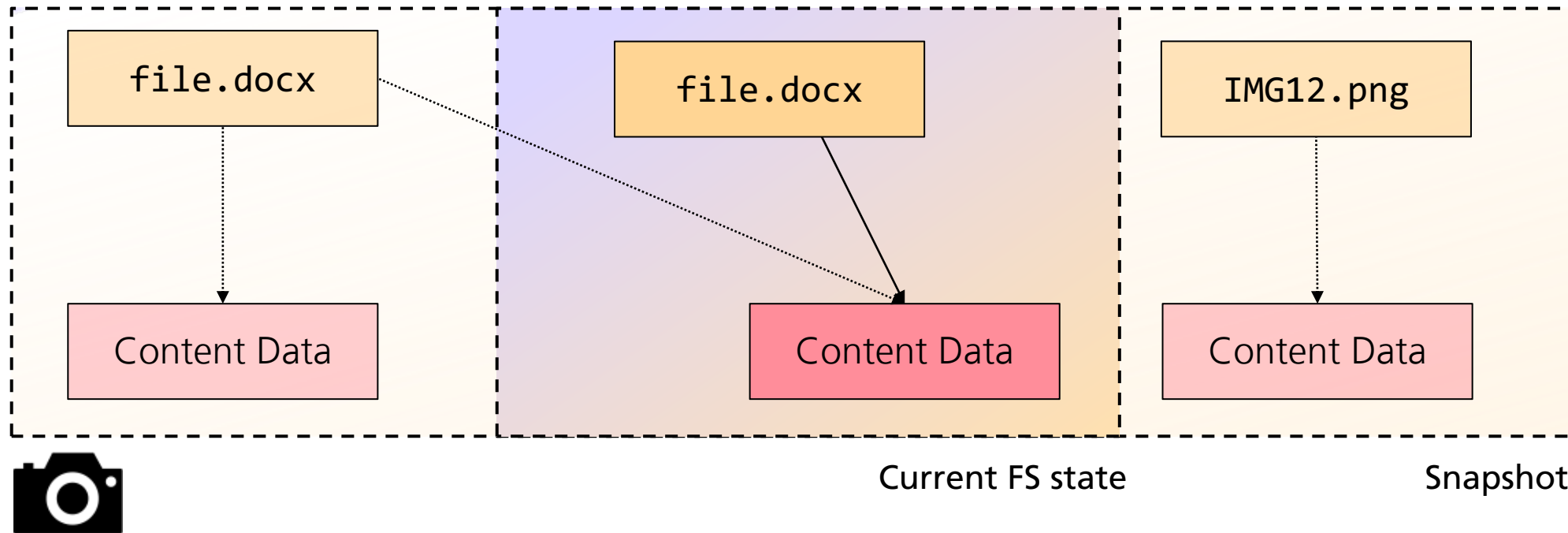
- Technique to ensure a file system stays in a **consistent state**
 - One possibility to ease the **creation of snapshots**



Copy-on-write

Possibility to go back in time!

- Technique to ensure a file system stays in a **consistent state**
 - One possibility to ease the **creation of snapshots**

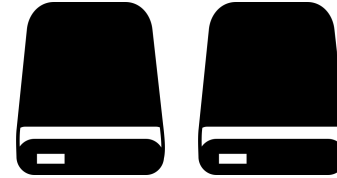




Fancy File System Techniques



Copy-on-write



Support for multiple
disks out of the box



Snapshots



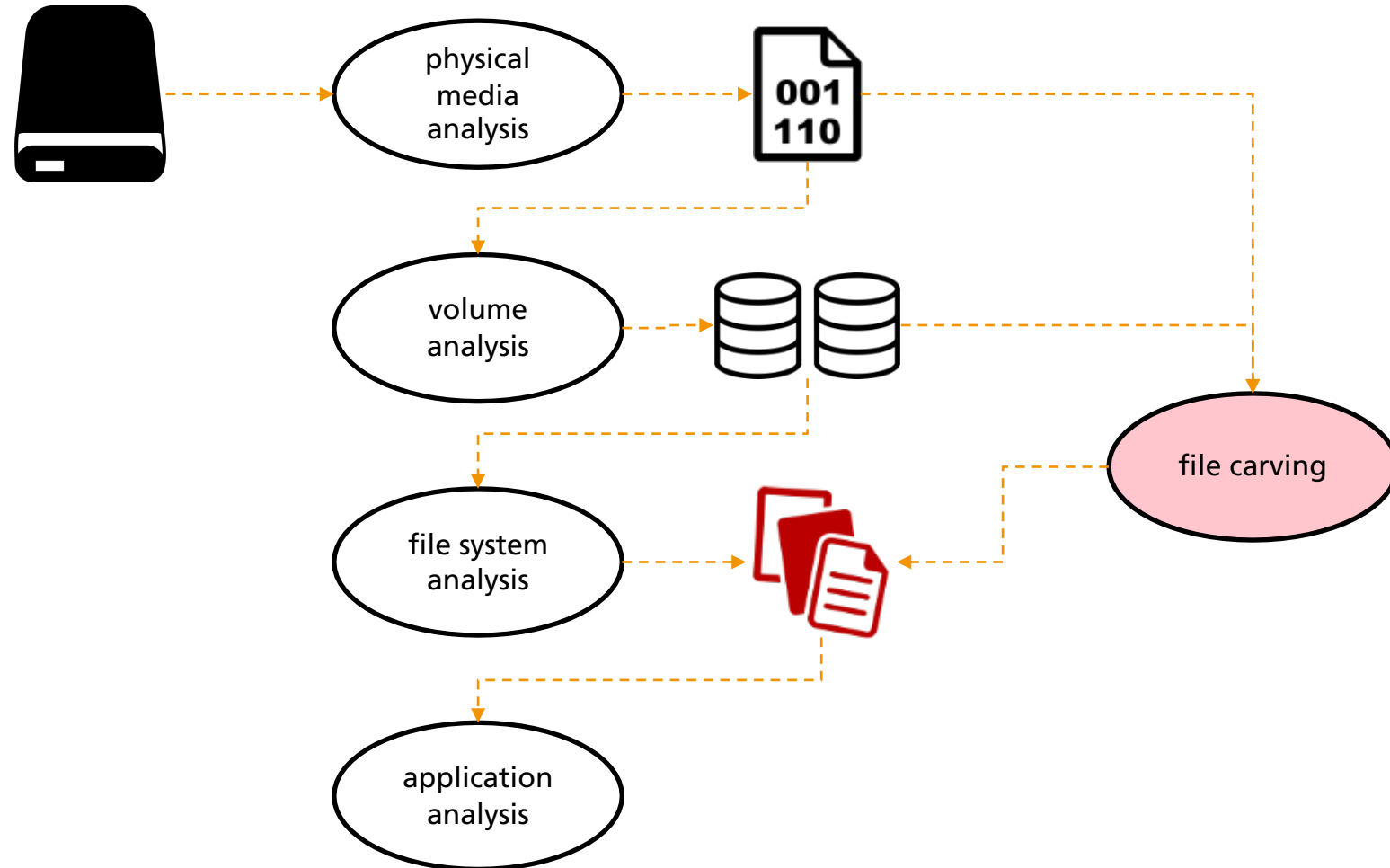
Encryption



Remote Conncetions

File System Analysis

What if we **cannot** use the file system in order to obtain our artifacts?





Any Questions?

